



SCELBI'S FIRST BOOK OF COMPUTER GAMES for the '8008/8080'

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Acknowledgement

Composing and typesetting material while working from rough drafts, and painstakingly copying source listings and the myriads of octal numbers from assembled listings, is a lot of tedious work which demands a high degree of accuracy. We wish to express our special appreciation for the services of:

Ms. Gabrielle Tingley

who performed this task for the material in this publication in a most efficient manner.

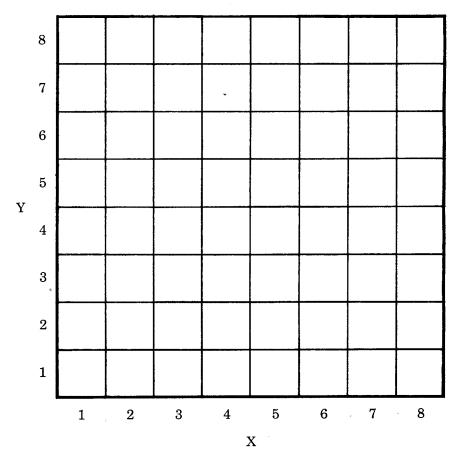
The Authors

SCELBI'S FIRST BOOK OF COMPUTER GAMES for the '8008/8080'

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SPACE CAPTURE

Space Capture is a game of skill and chance. The object of the game is to capture an imaginary space ship by destroying all the possible sectors that it might attempt to travel in. The game as presented herein utilizes a game board consisting of a grid containing 64 squares or sectors. The sectors are identified by X and Y coordinates. A pictorial of the playing board is illustrated below.

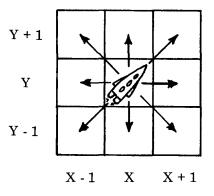


To start the game, the computer informs the player of a location

where the space ship was last observed. The player is then allowed to take one PHASOR SHOT by giving the X and Y coordinates of a SECTOR. The phasor shot will destroy the sector specified thereby preventing the space ship from traveling therein in the future. However, the player must be careful! If the space ship happens to be residing in a sector at the time it is destroyed by a phasor shot, then the space ship itself is considered destroyed. Since the object of the game is to CAPTURE the ship (for its cargo of course!), destroying the vessel is a losing move for the player.

The space ship is limited to moving only one sector at a time. As already mentioned, it may not move into an area that has been previously hit by a phasor shot. The ship's movement is also restricted to the boundaries of the eight by eight grid on which the game is played.

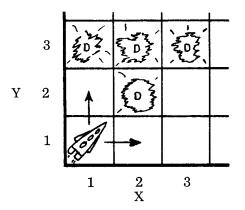
The maximum number of different moves the space ship has to choose from at any given time is thus eight. This is illustrated in the diagram shown below. This maximum number of possible moves, for instance, would be the case at the start of a game before a player had destroyed any sectors.



However, once the game is underway, the number of possible directions in which the space ship may move can be reduced. The example illustrated next shows the space ship in a position where only two moves are possible. This is because it is bounded on two sides by the edges of the playing grid. Additionally, the diagram

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shows several sectors marked by a D. These represent sectors that have been destroyed by the phasor shots of the player. The space ship may not enter into those areas. Thus, in the example, the space ship is only able to move up to the position X = 1 and Y = 2 or to the right into the position X = 2 and Y = 1. The space ship would be CAPTURED in the illustration if those two sectors had also been destroyed so that it could not move out of the indicated position X = 1 and Y = 1.



The game is relatively simple as far as computer games go, but it is a lot of fun because the moves of the space ship are made essentially random by the program. One may create strategies to attempt to use to capture the space ship, but one can never be certain where the next move will be until the space ship is captured. Also, if one does not take care where one shoots the phasor shots, the elements of chance can again enter the game. Remember, destroying a sector with the space ship traveling in it at the time ends the game!

The program for the game as it will be presented here will reside with room to spare in about 5 pages (256 bytes per page) of an '8008' or '8080' microcomputer system. If a person has even less memory available, the program can readily be compacted by the removal of some non-essential text messages. More room could be saved by more effective subroutining and attention to the program's organization to reduce the number of times pointers are altered. These techniques would allow the program to fit easily in less than three pages of memory. The reader should remember that the following program was designed so that the operation of the program could be easily followed. It was definitely not designed to minimize memory usage other than in the sense that this machine language version is many times more compact that would be required if the program utilized a higher level language for compilation or interpretation!

The fundamental operation of the program is outlined in the flow chart that appears on the next page. A brief verbal explanation of that chart will follow. Then the various portions of the program will be presented and discussed in detail.

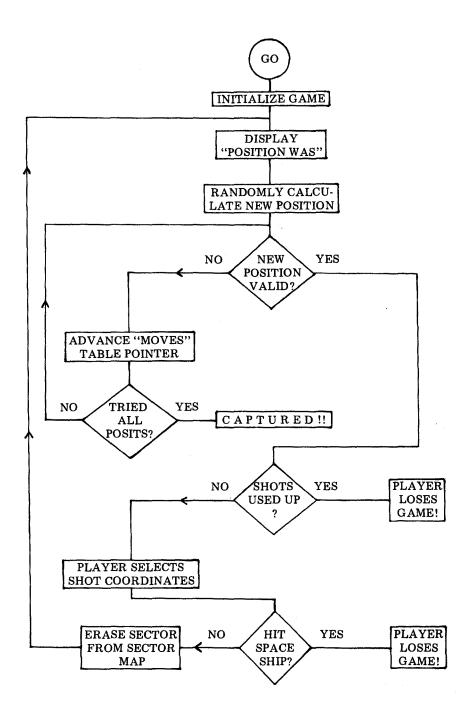
OUTLINE OF THE PROGRAM'S OPERATION

At the start of the program, a brief message is presented to explain the game to a new player. Next, the program determines if the operator desires to play a game. If not, a closing message is displayed and the program ends.

Assuming that a person elects to play a game, the program proceeds to select a semi-random starting point as the initial position of the space ship that is to be captured. The position of the space ship is then displayed as the WAS position to the operator. That is, the operator is informed of the LAST POSITION in which the space ship was observed. The operator then knows one position in which the space ship cannot be because the actual current position of the ship when the player fires a phasor shot will be in a sector adjacent to its last announced position. (Unless, of course, the ship has been captured.)

Once the WAS position has been displayed, the program proceeds to calculate a new position for the space ship to move into using an essentially random method. Whenever a calculation to move into a new sector has been made, the program must perform several tests. It must make sure that the new sector is within the bounds of the playing grid. And, it must make sure that it is not moving into a sector that has been destroyed by a player's phasor shot. If either of these tests fail, a new calculation is made to try

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another one of the eight possible adjacent sectors. If the program finds that all eight possible moves are blocked, the player wins. An appropriate message is then displayed.

Assuming that the space ship does have a valid move, the new position that it occupies is saved temporarily. If the player has not already expended an allotted number of shots, the program allows the operator to enter the coordinates of the sector that is to be eliminated from further occupation by the space ship. Once the coordinates have been entered, a test is made to see whether the space ship is presently in that sector. If so, the player loses as the space ship was destroyed versus being captured. If the space ship was not hit by the phasor shot, then the sector area is erased from a SECTOR MAP. Once an entry in the sector map has been erased, the space ship will be prevented from entering that sector in the future.

The game continues until the player uses up the allotted number of phasor shots, hits the space ship, or obtains a CAPTURE. At the conclusion of a game, the program queries the player as to whether a new game is to be played. Appropriate action is then taken as indicated above.

TEXT MESSAGES USED BY THE PROGRAM

Close to a third of the memory space utilized by the program is for storing the ASCII code for various messages that are displayed during the program's operation. These messages are of esthetic value particularly if the game is to be enjoyed by those who may not be familiar with the operation of a computer. The contents of these messages may be altered by the reader as desired, including complete deletion in many instances if one desires to conserve memory space. The various message strings that are used in the program being presented are listed next. "SPACESHIP CAPTURE. YOU HAVE 15 PHASOR SHOTS WITH WHICH TO DESTROY MY TRAVEL SECTORS. IF ALL MY ADJACENT SECTORS ARE DESTROYED I AM CAPTURED. IF YOU HIT ME OR RUN OUT OF PHASOR ENERGY, THEN YOU LOSE!"

"WANT TO PLAY?"

"POOR SPORT!"

"MY LAST POSITION WAS: X = "

", Y = "

"YOU ARE FIRING TO: X = "

YOU HIT ME!! YOU LOSE!"

YOU ARE OUT OF PHASOR ENERGY, YOU LOSE!"

"#!O# DARN! YOU HAVE ME CAPTURED!!"

The introductory message in particular takes almost a page of storage in memory and may readily be deleted if memory is at a premium in the user's system. The reader who wants to reduce the memory requirements some more can abbreviate the other messages if desired.

The text messages shown above are all stored in one continuous section in memory in the form of ASCII codes for the various characters in each string. (Note: In this manual the standard seven bit ASCII code will be shown with the code augmented by an eighth bit commonly referred to as the PARITY bit. The parity bit will always be assumed to be in the logic one or marking condition unless otherwise noted.)

A subroutine frequently referred to by the game program is shown

MSG,	LAM NDA RTZ	Fetch a character Set flags Finished if have zero byte
	CAL PRINT INL JFZ MSG INH	Else print character Advance low addr pointer Continue display Or adv page addr pointer
	JMP MSG	And then continue display

The MSG subroutine is quite simple. The calling program simply sets up the H and L memory pointing registers to the starting address of a string of characters that are to be outputted. Then, when the MSG subroutine is executed, the routine proceeds to fetch the characters from memory and output them until a zero byte is encountered. The subroutine itself calls on a subroutine labeled PRINT which must be provided by the program user. The PRINT subroutine must be an actual device operating routine that will cause the ASCII character in the accumulator to be transmitted to the output device being used by the system. The PRINT subroutine provided by the user may use the CPU registers B through E if required but it should not alter the contents of the H and L CPU registers. (Unless, of course, in doing so it is able to restore them to their original values before returning to the calling program.)

The reader will see that the MSG subroutine is used through-out the program being described. Prior to calling the subroutine, the main program will always setup the H and L registers to the starting address of the character string that is to be displayed. The character strings that will be used in the example program have been presented previously. It hardly goes without saying, that if a reader desires to modify the text messages, and by doing so alters their starting addresses, that appropriate modifications must be made to the setup address values whenever the MSG subroutine is used. This is also the case if the user decides to store the text messages at locations other than those shown in the program provided herein.

THE SPACE CAPTURE PROGRAM

The reader may refer to the flow chart presented earlier as the discussion of the actual operating portions of the program proceeds.

The first few procedures in the program consist of merely displaying the introductory message and then asking the prospective player if the playing of a game is desired.

Following the WANT TO PLAY query, the program then waits for a response from the system's input device which must be in the form of a letter Y for YES or N for NO. If a NO response is received at this point, then no game is to be played. The program will display the closing message and end the program. These first few operations are illustrated in the program listing below.

START,	LHI 000 LLI 000 CAL MSG	Pointer to introductory Message Display introductory message
OVER,	LHI 000	Pointer to WANT TO PLAY
	LLI 325	Message
	CAL MSG	Display message
INAGN,	CAL CKINP	See if have input
	NDA	Set flags
	CFS INPUTN	Fetch character if ready
	INL	Increment RANDOM cntr
	CPI 316	If input, was it N?
	JFZ NOTNO	Jump ahead if not N
	LHI 000	Pointer to POOR SPORT
	LLI 350	Message
	CAL MSG	Display message
	HLT	End of session

There are several instructions in the above routine that require elaboration. The reader may observe that there are references to two input subroutines. One is labeled CKINP. The other is designated

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INPUTN. The subroutine labeled INPUTN is a user created subroutine that will accept a character from the system's input device. Typically, this would be an ASCII encoded keyboard. The subroutine is expected to return the inputted character in the accumulator. This subroutine is free to use CPU registers B through E in performing its function. The INPUTN subroutine should also provide an echo capability by sending the character inputted out to the system's display device. This is done so the operator may verify the character inputted. (This might be accomplished by simply calling the previously mentioned PRINT subroutine.)

The CKINP subroutine is a user provided routine that simply performs a check to see if the input device has a character waiting to be inputted. If so, the subroutine must return with the MSB of the accumulator set to '0.' If a character is not ready, the subroutine should return with the most significant bit of the accumulator set to a logic '1' state.

The importance of having a separate subroutine (CKINP) that merely ascertains if a character is waiting will be explained here. The reader can see in the previous routine that if a character is ready to be received, the INPUTN subroutine will be called to actually obtain the data. However, whether or not a character is received, CPU register L will be incremented. CPU register L is actually used as a sort of random counter. The final value in register L will be determined by how long it takes for the player to respond with an input after a query from the program. This is because if there is no input the first time the instruction sequence is executed, the routine will eventually loop back to the point in the program labeled INAGN. Each time the program has to wait and goes back through the loop, the contents of register L are incremented. Naturally, this looping operation is being performed at a many-thousands-per-second rate.

How the final value in register L is used to form an essentially random number (when a valid input finally occurs) will be illustrated shortly. It is important to note that the inclusion of the separate CKINP subroutine is vital to the proper operation of the program being described. To reiterate, the CKINP subroutine must only ascertain if the input device has a character for the computer! It does not itself form a waiting loop for such a signal. That is accomplished by the previous routine in the manner described!

The program continues with a portion to be illustrated next starting with the label NOTNO. The first part of this sequence completes the test of the player's response to the WANT TO PLAY? query. This is done by testing to see if the character Y for YES was inputted. If not, the program loops back to the label INAGN just described to continue looking for a valid input.

When a Y response is received, the routine continues in the following manner. The value in register L is transferred to the accumulator. It is trimmed by a masking operation to leave only the three least significant bits. This would leave an octal number in the range of zero to seven. A count of one is added to this value to give a number in the range 01 to 10 octal. This is the equivalent of decimal 1 to 8, or the allowed coordinates along either the X or Y coordinate of the playing grid! This value is then saved in two temporary storage locations in memory to serve as the initial position of the space ship at the start of the game. Thus, the space ship will always have a starting point along a line corresponding to the diagonal where both coordinates are the same value. However, the actual value will vary with each game because of the random manner in which the number is generated (in register L) as alluded to previously. This gives some added variety to the game right from the beginning move!

The routine then continues by taking the value in the accumulator and reducing it by a masking operation back to the octal range 0 to 7. The value is then multiplied by 2 (RLC instruction) so that it will represent an even number in the range 00 to 16 octal.

At this point the value is converted to the low portion of an address. For this particular version of the program this is accomplished by the ORI 260 command which will form a value in the range 260 to 277 (octal). This address is stored temporarily in memory for use by a routine that will be explained in detail further on. Suffice it to say at this point that the address refers to a table that will contain the possible moves that the space ship might try to take.

The routine then continues by initializing a phasor shots taken counter that will keep track of how many shots the player has fired. Because of the point in the overall program at which the counter is decremented, this counter is initialized to a value one greater than the number of shots that the player is to be allowed.

The routine concludes by filling a block of 64 (decimal) locations with all ones. This block of memory will serve as a shots taken map. Its use will be explained in detail later.

NOTNO,	CPI 331 JFZ INAGN LAL NDI 007 ADI 001 LHI 001 LLI 372 LMA INL LMA LLI 377 NDI 007 RLC ORI 260 LMA LHI 001 LLI 376 LMI 020 LHI 003 LLI 300 LAI 377	If input, was it Y? If not, get a new input Else, move random counter To ACC & trim ASCII code Add 1 to get 1 - 8 Range and set up pointer To LAST position storage Initialize X WAS value Advance pointer Initialize Y WAS value Pointer to random cntr storage Reduce size Make it an even value Form table pointer And save table pointer Set pointer to shot counter Storage location Initialize to 16 decimal Set pointer to start of Shots taken map Fill accumulator with 1's
FILOOP,	LMA INL JFZ FILOOP	Initialize shots taken Map to all ones condition Until map completed

The next portion of the program starts by displaying the message MY LAST POSITION WAS: to the player. The routine then fetches the values of the X and Y coordinates that have been previously stored in memory and outputs those values by forming the ASCII code for the appropriate numerical values and displaying them via the output display subroutine PRINT provided by the user.

Next, a subroutine termed TRYMOV is called. The TRYMOV subroutine, which will be discussed shortly, will attempt to move the space ship into an available free sector using a technique that selects a new location in an essentially random manner. If the TRYMOV subroutine cannot move the space ship, the program will not return in the normal manner as the space ship will have been captured. If, however, the space ship is able to move to a new sector, the program will continue as illustrated in the routine. At this point, the phasor shots taken counter will be decremented in value. If the player has not used up the allotted shots, the game continues.

If the player has used up the number of allotted phasor shots, the program continues to the label PHASOR. Here the program will display the message indicating that the player is out of phasor energy and has lost the game. The program will then loop back to the label OVER presented earlier to see if the player wants to start a new game.

PLAYIN,	LHI 000	Set pntr to POSITION WAS:
	LLI 367	X = message
	CAL MSG	Display message
	LHI 001	Set pointer to X WAS
	LLI 372	Storage location
	LAM	Fetch value
	ORI 260	Form ASCII code
	CAL PRINT	Display position value
	LHI 001	Set pointer to Y =
	LLI 026	Message
	CAL MSG	Display message
	LHI 001	Set pointer to Y WAS
	LLI 373	Storage location
	LAM	Fetch value
	ORI 260	Form ASCII code
	CAL PRINT	Display position value
	CAL TRYMOV	Move the spaceship!

	LHI 001	Pointer to shots taken
	LLI 376	Counter storage
	LBM	Fetch counter
	DCB	Decrement value
	LMB	Restore counter
	JFZ CONTIN	Jump ahead if counter not 0
PHASOR,	LHI 001	If shots counter = 0 ,
	LLI 125	Set pointer and display
	CAL MSG	OUT OF ENERGY message
	JMP OVER	Go see if want new game

Provided that the player still has shots available with which to destroy travel sectors, the program continues at the label CONTIN. Here the message YOU ARE FIRING TO: is displayed. The program then allows the player to enter first the X and then the Y coordinate of the sector that the player wishes to destroy.

When obtaining the X coordinate, the program simply calls the subroutine INPUTN to obtain a character from the input device. The character obtained is checked to see if it is in the range of one to eight decimal. If not, the routine loops back to wait for a valid input. If so, the ASCII code is trimmed down to four bits and the value saved in a temporary location as the new value of X in memory.

The program then prepares to receive the Y coordinate from the player.

CONTIN,	LHI 001 LLI 036 CAL MSG	Set pointer to FIRING TO message Display message
INX,	CAL INPUTN CPI 261 JTS INX CPI 271 JFS INX LHI 001	Fetch X value See if input is a Digit in the range of one to eight decimal Ignore input if not If valid input set pointer

LLI 370	To new X storage location
NDI 017	Trim off ASCII part
LMA	Save the new X value
LHI 001	Set pointer to
LLI 026	Y = message
CAL MSG	Display message

The input for obtaining the Y coordinate from the player is handled in the same manner that was described for the portion of the program where the player responds to the WANT TO PLAY? query. Register L is again used as a counter whose final value will depend on how long it takes the player to enter the Y coordinate. When a valid character is received (in the range one to eight decimal), the trimmed number is saved in memory as the new coordinate along the Y axis. The value in register L is then processed in the same manner as before to form an address that will be utilized by the TRYMOV subroutine that has already been referred to, and which will be described soon.

INY,	CAL CKINP NDA CFS INPUTN INL CPI 261 JTS INY CPI 271 JFS INY LBL LHI 001 LLI 371 NDI 017 LMA LAB NDI 007 RLC ORI 260 LHI 001 LLI 377	See if have input Set flags Fetch character if ready Advance random counter See if input In decimal range One to eight Else ignore input Save random counter value Set pointer to new Y Storage location Trim ASCII part off Save the new Y value Move random counter to ACC Reduce it in size Make it an even value Form random table pointer Set pointer to random table Pointer storage location
	LLI 377 LMA	Pointer storage location Save random pointer

After the X and Y shot coordinates have been obtained from the player, the program continues at the point labeled HITEST. At this time the program must perform a check to determine whether the sector which the player has just destroyed is the same one in which the space ship might have been in. (Which is determined by the TRYMOV subroutine. Remember, the TRYMOV subroutine has already been called by the program, even though it has not yet been presented in detail in this discussion.) This is determined by testing to see if the coordinates of the player's phasor shot (stored in memory) match with the new location of the space ship (also stored in memory). If a match occurs here, then the space ship has been hit. That is a losing move for the player, and the program will display the appropriate message and return to see if the player wants to try a new game.

HITEST,	LHI 001 LLI 370 LAM INL INL CPM JFZ ZERSEC DCL LAM INL INL CPM JFZ ZERSEC	Set pointer to X phasor shot storage Fetch X shot value Advance pointer to new X spaceship location value Compare shot with location If not a match, no hit Set pointer to Y phasor Shot storage and fetch Advance pointer to new Y spaceship location value Compare shot with location If not a match, no hit
BOMB,	LHI 001 LLI 072 CAL MSG JMP OVER	Shot hit spaceship - set Pointer to HIT message Display message See if want new game

The next section of the program begins at the label ZERSEC. This portion of the program serves to zero-out a sector in the sector map whenever a sector is destroyed by a phasor shot made by the player.

As the reader knows, there are 64 different sectors in the 8 by 8

grid on which the game is played. At the beginning of a game, an area in memory is assigned as a sector map. This area consists of 64 consecutive locations in memory. In the sample program, the memory area is assigned to locations 300 to 377 (octal) on page 03. The area is initialized (as mentioned earlier) by loading the value 377 into all 64 locations. Now, each time the player specifies a grid location by designating an X and Y coordinate, a memory location in the sector map is changed to be in a 000 (octal) condition. The location that is to be zeroed out is ascertained by performing a simple calculation. The fundamental calculation made may be expressed by the following formula:

$$V = [(X - 1) \times 8] + (Y - 1)$$

where X and Y are the respective coordinates given by the player when firing a phasor shot. The value V obtained by the calculation is then added to a base value (300 in the example program) to give an effective address in the sector map. By reviewing the above formula, the reader may verify that the calculation will yield the values from 0 to 63 decimal or 0 to 77 octal when all the possible coordinate values are considered. When added to the base value (300 in the example) this will yield a low address in the range 300 to 377 (octal).

A portion of the program (to be presented later) will prevent the space ship from moving into any location that has been zeroed-out in the sector map.

The ZERSEC portion of the program is presented next.

ZERSEC,	LLI 370 LAM	Get X shot value From storage
	SUI 001	Subtract '1'
	RLC	
	RLC	Multiply by eight
	RLC	
	LDA	Save in register D temporarily
	INL	Get Y shot value
	LAM	From storage
	RLC RLC LDA INL	Save in register D temporarily Get Y shot value

SUI 001	Subtract '1'
ADD	Add to previous calculations
ORI 300	Form shot table address
LLA	Set low address pointer
LHI 003	And page address of shot table
LMI 000	Zero the entry in shot table
JMP PLAYIN	Continue with game

The next portion of the program is the subroutine TRYMOV. This subroutine is the most complicated portion of the program. The subroutine serves the function of attempting to find a new location for the space ship. It does this by attempting to find a sector adjacent to the last position of the space ship that has not been destroyed by a phasor shot. The sector must also be within the boundaries of the 8 by 8 playing grid. Additionally, the direction of movement is accomplished in an essentially random manner so that the player will not be able to detect a reliable pattern of movement for the space ship!

The first few instructions of the subroutine fetch the address that points to the move table. This address is set up each time the player specifies the Y coordinate of a phasor shot (or answers the WANT TO PLAY? query at the beginning of a game). As discussed earlier, the address each time the subroutine is entered will have been selected in an essentially random manner.

The address refers to a position in a table which holds all the possible moves the space ship can make to an adjacent sector. A pictorial near the beginning of this article illustrated the eight possible moves the space ship could make if it was not bounded by the edges of the playing grid, or sectors that had been destroyed. Referral to that pictorial will show that the possible moves may be referenced by a value of -1, 0, or +1 from its present position along each axis. One can convert this information into a table that holds all the possible moves. The table used consists of eight groups of two bytes per group. The first byte in a group holds a move along the X axis. The second stores a move along the Y axis. In the example program, the table is stored in locations 260 through 277 on page 03, and appears as shown on the following page.

377	MOVES TABLE $(X = -1)$
001	Y = +1
000	$\mathbf{X} = 0$
001	Y = +1
001	X = +1
001	Y = +1
377	X = -1
000	$\mathbf{Y} = 0$
001	X = +1
000	$\mathbf{Y} = 0$
377	X = -1
377	Y = -1
000	$\mathbf{X} = 0$
377	Y = -1
001	X = +1
377	Y = -1

The initial value of the address to the moves table is transferred from a temporary location in memory to the accumulator at the start of the TRYMOV subroutine. Also, a moves tried counter, which will be maintained in CPU register C during the subroutine, is initialized to a value of eight (decimal). The routine then goes on to the point labeled TRYSEC which marks a looping point within the subroutine.

At TRYSEC, the address value originally in the accumulator is moved to CPU register L to set up the low portion of the address to the moves table. This value is also saved for possible later use in CPU register B. The high address of the moves table is set up in CPU register H, and an X move value fetched from the table. The X move value obtained from the table is then added to the X coordinate value representing the previous position of the space ship to form a new value along the X coordinate. At this point some tests must be made to ensure that the new value is within the boundaries of the playing grid. This is readily accomplished by checking to see that the new coordinate is between the range of 1 to 8 decimal. If the new value is within the playing grid, it is saved in a temporary location in memory. If it is not valid, the program jumps ahead to a routine that will advance the moves table pointer to the next X entry in the table (by going to the label NOGDX which will advance the address temporarily stored in CPU register B TWO locations!)

If the new X value is O.K., the routine proceeds to advance the pointer to the moves table one location to obtain a Y move. A similar boundary checking procedure is performed again. If the new Y coordinate is valid, it too is saved in a temporary location in memory. If not, the program jumps ahead to the label NOGDY which will advance the moves table pointer just one location to the next X entry.

If the new X and Y coordinates are within the boundaries of the playing grid, the program continues by executing the portion of the subroutine labeled CHECK. This part of the program ascertains whether the new sector the space ship is attempting to move into is available. That is, that it has not been destroyed by a previous phasor shot made by the player! This is readily accomplished by using the new coordinate values to once again calculate a position in the sector map and checking to see that the position has not been zeroed out. The calculation technique is exactly the same as that used by the routine ZERSEC explained earlier. If the sector is available, the program jumps ahead to the routine labeled SAVPOS. If not, the space ship must try to find another position in the grid. This is attempted by proceeding to the point labeled NOGDY.

The portion of the routine beginning with the label NOGDX serves to advance the moves table address pointer to the next X entry in the table. Since the table occupies just 20 (octal) locations in memory (in the example program locations 260 through 277 on page 03), and since the initial value may have been at any even valued address within that range, some special operations must be performed when advancing the pointer to the next X entry. First, it may be necessary to try every possible move in the table. Since the first position tried may have been the eighth entry in the table (four least significant bits of the address equal to 16 octal), one must keep the pointer in the range of 00 to 17 octal (for the four LSB's). This is readily accomplished by a masking operation that removes the four most significant bits. Then, since the table does not reside in the first 20 (octal) locations on a page in memory, the base address value of 260 (in the example) must be tacked back on to form the complete address value. This procedure will force the moves table pointer to loop back to the value 260 when it reaches a value of 277 rather than going to 300 which would be outside the range of the table.

When the address of the next entry in the moves table has been set up, the routine checks to see if all eight possible moves have been tried by decrementing the counter being maintained in CPU register B. If not, the routine loops back to the point labeled TRYSEC. If, however, all eight possible moves have been tried, then the space ship has not been able to find a new position and it is captured. A message of defeat is then issued to the player. From that point, the program will go back to see if a new game is desired.

The portion of the routine labeled SAVPOS is executed if the space ship successfully finds a new sector to move into. This routine simply moves the X and Y coordinate values being temporarily saved in memory during the TRYMOV subroutine into the storage locations used to hold the new location. (This move actually places the new coordinates into the memory locations that are referenced the next time the message MY LAST LOCATION WAS: is displayed.)

Note that the SAVPOS routine ends with a RET instruction to conclude the subroutine. Thus, a return from the TRYMOV subroutine indicates that the space ship completed a successful move. If the space ship does not find a new sector to move into, the subroutine is not actually completed. Instead, a jump out of the subroutine to start a new game is executed.

The various portions of the TRYMOV subroutine are presented below.

TRYMOV,	LHI 001 LLI 377 LAM LCI 010	Set pointer to Random counter storage Fetch value Set a loop counter
TRYSEC,	LLA LBA LHI 003	Set pointer to moves table Save pointer in B too Set pg pointer to moves table

LAM	Fetch an X move
LHI 001	Change pointer to
LLI 372	X WAS storage
ADM	Add X move to form new loc
CPI 001	Now make boundaries test
JTS NOGDX	No good if less than one
CPI 011	Or more than
JFS NOGDX	Eight decimal
LHI 001	0
LLI 374	If OK, save in X temporarily
LLI 574 LMA	Storage location For awhile
INB	Advance pointer stored in B
LLB	And load new pointer
LHI 003	To Y move location
	Fetch a Y move
LHI 001	Change pointer to
LLI 373	Y WAS storage
ADM	Add Y move to form new loc
CPI 001	Now make boundaries test
JTS NOGDY	No good if less than one
CPI 011	Or more than
JFS NOGDY	Eight decimal
LHI 001	If OK, save in Y temporarily
LLI 375	Storage location
LMA	For awhile
DCL	Decrement pointer back to
LAM	X temp storage and fetch
SUI 001	Subtract '1'
RLC	
RLC	Multiply by eight
RLC	
LDA	Save in D temporarily
INL	Advance pointer to
LAM	Y temp storage and fetch
SUI 001	Subtract '1'
ADD	Add to previous calculations
ORI 300	Form shot table address
LLA	Set low address pointer
LHI 003	And page address pointer
	T G T C P C P T C P P -

CHECK,

	LAM NDA JFZ SAVPOS JMP NOGDY	Fetch entry from shot table Set flags, now see if location Previously fired into Try another location if yes
NOGDX, NOGDY,	INB INB LAB NDI 017 ORI 260 DCC JFZ TRYSEC LHI 001 LLI 201 CAL MSG JMP OVER	Advance move table pointer As required to get to next X move, move value to ACC And make sure it is kept In bounds Decrement loop counter If not 0, try another location Else set pointer To CAPTURED message Display message See if want new game
SAVPOS,	LHI 001 LLI 374 LDM INL LEM LLI 372 LMD INL LME RET	Set pointer to X temporarily Storage location Save value in D temporarily Advance pointer to Y temp Save in E temporarily Change pointer to X WAS Storage and set new value Do likewise for Y WAS Too! Return to calling program

That is all there is to the program! That is not so difficult, eh?

An assembled listing of the program for running on a 8008 system will be presented on the following pages. The program in the listing has been assembled to reside in pages 01 to 03 with page 04 reserved for the user's I/O routines. Page 00 and most of page 01 will be used to hold the various message strings mentioned previously. The ASCII data that should be stored in those locations is shown next. (Assuming the user is satisfied with the message strings illustrated.)

000	000	215	212	212	323	320	301	303	305
000	010	323	310	311	320	240	303	301	320
000	020	324	325	322	305	256	240	331	317
000	030	325	240	310	301	326	305	240	261
000	040	265	240	320	310	301	323	317	322
000	050	215	212	323	310	317	324	323	240
000	060	327	311	324	310	240	327	310	311
000	070	303	310	240	324	317	240	304	305
000	100	323	324	322	317	331	240	315	331
000	110	240	324	322	301	326	305	314	215
000	120	212	323	305	303	324	317	322	323
000	130	256	240	311	306	240	301	314	314
000	140	240	315	331	240	301	304	312	301
000	150	303	305	316	324	240	323	305	303
000	160	324	317	322	323	215	212	301	322
000	170	305	240	304	305	323	324	322	317
000	200	331	305	304	240	311	240	301	315
000	210	240	303	301	320	324	325	322	305
000	220	304	256	240	311	306	240	331	317
000	230	325	215	212	310	311	324	240	315
000	240	305	240	317	322	240	322	325	316
000	250	240	317	325	324	240	317	306	240
000	260	320	310	301	323	317	322	240	305
000	270	316	305	322	307	331	254	215	212
000	300	324	310	305	316	240	331	317	325
000	310	240	314	317	323	305	241	215	212
000	320	212	000	000	000	000	215	212	212
000	330	327	301	316	324	240	324	317	240
000	340	320	314	301	331	277	240	240	000
000	350	215	212	212	320	317	317	322	240
000	360	323	320	317	322	324	241	000	215
000	370	212	212	315	331	240	314	301	323
001	000	324	240	320	317	323	311	324	311
001	010	317	316	240	327	301	323	272	240
001	020	240	330	240	275	240	000	254	240
001	030	240	331	240	275	240	000	215	212
001	040	212	331	317	325	240	301	322	305
001	050	240	306	311	322	311	316	307	240
001	060	324	317	272	240	240	330	240	275
001	070	240	000	215	212	212	331	317	325

1 - 24

$\begin{array}{ccc} 001 & 100 \\ 001 & 110 \\ 001 & 120 \end{array}$	$\begin{array}{ccc} 240 & 310 \\ 241 & 240 \\ 317 & 323 \end{array}$	240	331	317	325	240	314
$\begin{array}{c} 001 \ 130 \\ 001 \ 140 \end{array}$	$331 \ 31' \ 317 \ 32!$						
$\begin{array}{ccc} 001 & 150 \\ 001 & 160 \end{array}$	$310 \ 301 \ 305 \ 325$						
$\begin{array}{ccc} 001 & 170 \\ 001 & 200 \end{array}$	$317 \ 320 \ 000 \ 210$						
$\begin{array}{ccc} 001 & 210 \\ 001 & 220 \end{array}$	$\begin{array}{c} 207 & \overline{20'} \\ 316 & 24 \end{array}$				-		-
$\begin{array}{c} 001 \ \ 230 \\ 001 \ \ 240 \end{array}$	$310 \ 302$ $240 \ 302$						
$\begin{array}{ccc} 001 & 250 \\ 001 & 260 \end{array}$	$240 \ 322$ $240 \ 242$			240	305	2 40	304

The starting address for each message string may be located from the above data presentation. The beginning of each message string has been underlined. The reader may desire to change some of the messages. If the reader elects to do so, and by so doing changes the starting address of a character string, then the appropriate pointer instruction in the operating portion of the program must be modified accordingly. The assembled program for an 8008 system is presented on the next several pages.

001	350	307			MSG,	LAM
001	351	240				NDA
001	352	053				RTZ
001	353	106	200	004		CAL PRINT
001	356	060				INL
001	357	110	350	001		JFZ MSG
001	362	050				INŲ
001	363	104	350	001		JMP MSG
001	370	000				000
001	371	000				000
001	372	000				000
001	373	000				000

001	374	000				000
001	375	000				000
001	376	000				000
001	377	000				000
002	000	056	000		START,	LHI 000
002	002	066	000			LLI 000
002	004	106	350	001		CAL MSG
002	007	056	000		OVER,	LHI 000
002	011	066	325			LLI 325
002	013	106	350	001		CAL MSG
002	016	106	000	004	INAGN,	CAL CKINP
002	021	240				NDA
002	022	122	020	004		CFS INPUTN
002	025	060				INL
002	026	074	316			CPI 316
002	030	110	043	002		JFZ NOTNO
002	033	056	000			LHI 000
002	035	066	350			LLI 350
002	037	106	350	001		CAL MSG
002	042	000				HLT
002	043	074	331		NOTNO,	CPI 331
002	045	110	016	002		JFZ INAGN
002	050	306				LAL
002	051	044	007			NDI 007
002	053	004	001			ADI 001
002	055	056	001			LHI 001
002	057	066	372			LLI 372
002	061	370				LMA
002	062	060				INL
002	063	370				LMA
002	064	066	377			LLI 377
002	066		007			NDI 007
002	070	002				RLC
	071		260			ORI 260
	073	370				LMA
	074		001			LHI 001

002	076	066	376			LLI 376
	100		020			LMI 020
002	102	056	003			LHI 003
002	104	066	300			LLI 300
002	106	006	377			LAI 377
002	110	370			FILOOP,	LMA
002	111	060			,	INL
	112	110	110	002		JFZ FILOOP
002	115	056	000		PLAYIN,	LHI 000
002	117	066	367			LLI 367
002	121	106	350	001		CAL MSG
002	124	056	001			LHI 001
002	126	066	372			LLI 372
002	130	307				LAM
002	131	064	260			ORI 260
002	133	106	200	004		CAL PRINT
002	136	056	001			LHI 001
002	140	066	026			LLI 026
002	142	106	350	001		CAL MSG
002	145	056	001			LHI 001
002	147	066	373			LLI 373
002	151	307				LAM
002	152	064	260			ORI 260
002	154	106	200	004		CAL PRINT
002	157	106	002	003		CAL TRYMOV
002	162	056	001			LHI 001
002	164	066	376			LLI 376
002	166	317				LBM
002	167	011				DCB
002	170	371				LMB
002	171	110	206	002		JFZ CONTIN
002	174	056	001		PHASOR,	LHI 001
002	176	066	125			LLI 125
002	200	106	350	001		CAL MSG
	203		007			JMP OVER
002	206	056	001		CONTIN,	LHI 001
		-	_		,	

	$\begin{array}{c} 210\\ 212 \end{array}$		036 350	001		LLI 036 CAL MSG
002	$\begin{array}{c} 215 \\ 220 \end{array}$	074			INX,	CAL INPUTN CPI 261
	222		215	002		JTS INX
	225	074				CPI 271
	227		215	002		JFS INX
	232	056				LHI 001
	234	066				LLI 370
	236	044	017			NDI 017
002	240	370				LMA
002	241	056				LHI 001
	243					LLI 026
	245	106	350	001		CAL MSG
002	250	106	000	004	INY.	CAL CKINP
002	253	240			7	NDA
002	254		020	004		CFS INPUTN
002	257	060				INL
002	260	074	261			CPI 261
	262	160	250	002		JTS INY
002	265 267 272	074	271			CPI 271
002	267	120	250	002		JFS INY
002	272	316				LBL
002	273	056	001			LHI 001
002	275	066	371			LLI 371
	277	044	017			NDI 017
002		370				LMA
	302	301				LAB
	303	044	007			NDI 007
002	305	002				RLC
002	306	064	260			ORI 260
002	310	056	001			LHI 001
002	312	066	377			LLI 377
002	314	370				LMA
	315	000			HITEST,	LHI 001
002		066	370		,	LLI 370
002	321	307				LAM

002 002 002 002 002 002 002 002	325 330 331 332 333	061 307 060 060 277	352 352			INL INL CPM JFZ ZERSEC DCL LAM INL INL INL CPM JFZ ZERSEC
002	340	056			BOMB,	LHI 001
002		066				LLI 072
002	344		350			CAL MSG
002	347	104	007	002		JMP OVER
	352	066	370		ZERSEC,	LLI 370
	354	307				LAM
002		024	001			SUI 001
002		002				RLC
002		002				RLC
002		002				RLC
002		330				LDA
002		060				INL
	354	307				LAM
002		024	001			SUI 001
002		203				ADD
	370	064	300			ORI 300
	372	360				LLA
	373	056	003			LHI 003
	375	076	000			LMI 000
002	377	104	115	002		JMP PLAYIN
003		056			TRYMOV,	LHI 001
	004	066				LLI 377
	006	307				LAM
003	007	026	010			LCI 010
	011				TRYSEC,	LLA
003	012	310				LBA

003	013	056	003			LHI 003
003		307	000			LAM
003		056	001			LHI 001
003		066				LLI 372
003		207	012			ADM
003		074	001			CPI 001
003			125	003		JTS NOGDX
003		074		000		CPI 011
003			125	003		JFS NOGDX
003		056		000		LHI 001
003		066				LLI 374
003		370	0.1			LMA
003		010				INB
003		361				LLB
	044	056	003			LHI 003
003	046	307				LAM
003	047	056	001			LHI 001
003	051	066	373			LLI 373
003	053	207				ADM
003	054	074	001			CPI 001
003	056	160	126	003		JTS NOGDY
003	061	074	011			CPI 011
003	063	120	126	003		JFS NOGDY
003	066	056	001			LHI 001
	070	066	375			LLI 375
003	072	370				LMA
	073	061			CHECK,	DCL
	074	307				LAM
	075	024	001			SUI 001
003		002				RLC
003		002				RLC
003		002				RLC
003		330				LDA
	103	060				INL
	104	307	001			LAM
	105	024	001			SUI 001
	107	203	000			ADD
003		064	300			ORI 300
003	112	360				LLA

003 11 003 11 003 11 003 11 003 12	15 307 16 240 17 110	003 152 126			LHI 003 LAM NDA JFZ SAVPOS JMP NOGDY
003 12 003 12 003 12 003 13 003 13 003 13 003 14 003 14 003 14 003 14 003 14 003 14 003 14 003 14 003 14 003 14 003 14	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	017 260 011 001 201 350 007	001	NOGDX, NOGDY,	
003 18 003 18 003 18 003 18 003 16 003 16 003 16 003 16	52 056 54 066 56 337 57 060 50 347 51 066 53 373 54 060 55 374	001 374 372		SAVPOS,	
003 26 003 26 003 26 003 26 003 26 003 26 003 26 003 27 003 27	31 001 32 000 33 001 34 001 35 001 36 377 37 000 70 001				377 001 000 001 001 001 377 000 001 000

003 002 003 003 003 003	273 274 275 276	377 377 000 377 001 377	377 377 000 377 001 377
004	000	CKIN	NP,
004	020	INPUT	'N,
004	200	PRIN	NT,

Do not forget that the program as presented will be using locations 300 through 377 on page 03 as a sector map. The reader should also make sure that the user provided I/O routines are loaded into memory at the indicated locations before attempting to operate the program!

OPERATING THE SPACE CAPTURE PROGRAM

Once the program has been loaded into memory it is ready for operation. The program is started by executing a jump to location 000 on page 02 for the illustrated program, and placing the computer in the normal program execution run mode. From there on the program effectively guides the player. The program will continue to operate, playing game after game, until the player responds with a N for NO to the WANT TO PLAY? query.

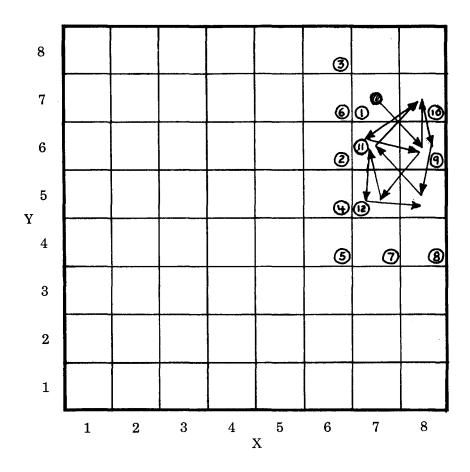
The player will want to have a supply of paper with 8 by 8 grids marked out to keep track of the space ship's movements and sectors in which shots have been fired as the games progresses. If the game is to be used frequently, it is probably worthwhile to make up a good supply of the grid forms using a mimeograph or duplicating machine.

In case the reader has any doubts as to how the game is played,

the following illustrates an actual game played using the program. At the end of the illustration showing the dialogue between the computer and the player is a grid illustrating how the game progressed. The progress of the space ship is shown as a series of arrows indicating the direction of each movement. The phasor shots fired by the player are shown as a circled number in various sectors. The number refers to the actual shot number as the game progressed.

> SPACESHIP CAPTURE. YOU HAVE 15 PHASOR SHOTS WITH WHICH TO DESTROY MY TRAVEL SECTORS. IF ALL MY ADJACENT SECTORS ARE DESTROYED, I AM CAPTURED. IF YOU HIT ME OR RUN OUT OF PHASOR ENERGY, THEN YOU LOSE!

WANT TO PLAY? Y MY LAST POSITION WAS: X=7, Y=7YOU ARE FIRING TO: X=7, Y=7MY LAST POSITION WAS: X=8, Y=6YOU ARE FIRING TO: X=6, Y=6MY LAST POSITION WAS: X=8, Y=7YOU ARE FIRING TO: X=6, Y=8MY LAST POSITION WAS: X=8, Y=6YOU ARE FIRING TO: X=6, Y=5MY LAST POSITION WAS: X=8, Y=5YOU ARE FIRING TO: X=6, Y=4MY LAST POSITION WAS: X=7, Y=6 YOU ARE FIRING TO: X = 6, Y = 7MY LAST POSITION WAS: X = 8, Y = 7YOU ARE FIRING TO: X = 7, Y = 4MY LAST POSITION WAS: X = 7, Y = 6YOU ARE FIRING TO: X = 8, Y = 4MY LAST POSITION WAS: X = 8, Y = 6YOU ARE FIRING TO: X = 8, Y = 6MY LAST POSITION WAS: X = 7, Y = 5YOU ARE FIRING TO: X = 8, Y = 7MY LAST POSITION WAS: X = 7, Y = 6YOU ARE FIRING TO: X = 7, Y = 6MY LAST POSITION WAS: X = 7, Y = 5YOU ARE FIRING TO: X = 7, Y = 5MY LAST POSITION WAS: X = 8, Y = 5**#**!0**#**! DARN! YOU HAVE ME CAPTURED!! WANT TO PLAY?



PICTORIAL OF THE MOVES MADE IN THE ILLUSTRATIVE SPACE CAPTURE GAME

LISTING FOR AN 8080 COMPUTER

The following is a listing of the program for an 8080 system. Only a few minor changes have been made in the program. Notably, the inclusion of stack pointer initializing instructions (required by the 8080 since it does not have a program counter stack on the CPU chip) at the labels START and OVER. Additionally, the double register (H and L) load instruction has been utilized when applicable instead of the individual commands required in an 8008 unit. Several other minor changes have been made to make use of the more powerful 8080 instruction set, but the basic structure of the program has not been altered so that the explanations of the various routines made earlier need not be elaborated upon.

001 001 001 001	350 351 352 353 356 357	043	200 350		MSG,	LAM NDA RTZ CAL PRINT INXH JMP MSG
	370	000				000
001	371	000				000
001	372	000				000
001	373	000				000
001	374	000				000
001	375	000				000
001	376	000				000
001	377	000				000
002	000	061	350	001	START,	LXS 350 001
	003		000		~	LXH 000 000
	006		350			CAL MSG
002	011	061	350	001	OVER,	LXS 350 001
002	014	041	325	000		LXH 325 000
002	017	315	350	001		CAL MSG

002 002 002 002 002 002		247 364 054 376 302 041	000 020 316 046 350 350	004 002 000	INAGN,	CAL CKINP NDA CFS INPUTN INL CPI 316 JFZ NOTNO LXH 350 000 CAL MSG
002	045	166				HLT
002 002 002	046 050 053 054 056			002	NOTNO,	CPI 331 JFZ INAGN LAL NDI 007 ADI 001
002	060	041	372	001		LXH 372 001
	063	167				LMA
		054				INL
	065	167	077			LMA
	066 070	$\begin{array}{c} 056 \\ 346 \end{array}$				LLI 377
	070	007	007			NDI 007 RLC
	072	366	260			ORI 260
	075	167	200			LMA
	076		376	001		LXH 376 001
	101		020			LMI 020
002	103	041	300	003		LXH 300 003
002	106	076	377			LAI 377
002	110	167			FILOOP,	LMA
002	111	054			,	INL
002	112	302	110	002		JFZ FILOOP
	115		367		PLAYIN,	LXH 367 000
	120		350			CAL MSG
	123		372	001		LXH 372 001
	126	176				LAM
	127		260			ORI 260
002	131	315	200	004		CAL PRINT

002			026	
002	137	315	350	001
	142	041	373	001
002	145	176		
002	146	366	260	
002	150	315	200	004
002	153	315	363	002
002		041		
002		065		
002		302	176	002
002	102	002	110	002
002	165	041	125	001
002		315		
		303		
002	173	303	011	002
	1 = 0	0.41	000	0.01
	176	041	036	
002	201	315	350	001
	204		020	004
	207		261	
	211		204	002
002	214	376	271	
002	216	362	204	002
002	221	041	370	001
002	224	346	017	
002	226	167		
002	227	041	026	001
002		315	350	001
002	235	315	000	004
	240	247		
	241	364	020	004
	244	054	040	001
	245	376	261	
002	243 247	372		002
002				002
	$\begin{array}{c} 252 \\ 254 \end{array}$	376		000
		362	235	002
	257	105	0	0.07
002		041	371	001
002	263	346	017	

	LXH 026 001 CAL MSG LXH 373 001 LAM ORI 260 CAL PRINT CAL TRYMOV LXH 376 001 DCM JFZ CONTIN
PHASOR,	LXH 125 001 CAL MSG JMP OVER
CONTIN,	LXH 036 001 CAL MSG
INX,	CAL INPUTN CPI 261 JTS INX CPI 271 JFS INX LXH 370 001 NDI 017 LMA LXH 026 001 CAL MSG
INY,	CAL CKINP NDA CFS INPUTN INL CPI 261 JTS INY CPI 271 JFS INY LBL LXH 371 001 NDI 017

002 002 002 002	265 266 267 271 272 274 277	366	260	001		LMA LAB NDI 007 RLC ORI 260 LXH 377 001 LMA
002 002 002 002 002 002 002	300 303 304 305 306 307 312 313	176 054 054		001	HITEST,	LXH 370 001 LAM INL INL CPM JFZ ZERSEC DCL LAM
002 002 002 002	314 315 316	054 054 276 302	333 072		BOMB	INL INL CPM JFZ ZERSEC LXH 072 001
002	325 330	315	350 011	001	DOMD,	CAL MSG JMP OVER
000	000	050	050		775	T T T O T O
002 002 002 002 002 002 002 002 002 002	343 344 345 346 350 351 353 354	$\begin{array}{c} 056\\ 176\\ 326\\ 007\\ 007\\ 127\\ 054\\ 176\\ 326\\ 202\\ 366\\ 157\\ 046\\ 066\\ \end{array}$	001 001 300 003		ZERSEC,	LLI 370 LAM SUI 001 RLC RLC LDA INL LAM SUI 001 ADD ORI 300 LLA LHI 003 LMI 000

002	360	303	115	002		JMP PLAYIN
$\begin{array}{c} 002 \\ 002 \\ 002 \end{array}$		041 176 016	377 010	001	TRYMOV,	LXH 377 001 LAM LCI 010
002 003 003 003 003 003 003 003 003 003	$\begin{array}{c} 372\\ 373\\ 375\\ 376\\ 001\\ 002\\ 004\\ 007\\ 011\\ 014\\ 017\\ 020\\ 021\\ 022\\ 024\\ 025\\ 030\\ 031\\ 033\\ 036\\ 040 \end{array}$	$\begin{array}{c} 206\\ 376\\ 372\\ 376\\ 362\\ 041\\ 167\\ 004\\ 150\\ 046\\ 176\\ 041\\ 206\\ 376\\ 372\\ 376\\ 362\\ \end{array}$	 372 001 101 011 101 374 003 373 001 102 	003 003 001 001 003 003	TRYSEC,	LLA LBA LHI 003 LAM LXH 372 001 ADM CPI 001 JTS NOGDX CPI 011 JFS NOGDX LXH 374 001 LMA INB LLB LHI 003 LAM LXH 373 001 ADM CPI 001 JTS NOGDY CPI 011 JFS NOGDY LXH 375 001
003		167	010	001		LMA
003 003 003 003 003 003 003 003 003	050 051 053 054 055 056 057	055 176 326 007 007 127 054 176	001		CHECK,	DCL LAM SUI 001 RLC RLC RLC LDA INL LAM

003 003 003 003 003 003 003 003 003	063 064 066 067 071 072 073		300			SUI 001 ADD ORI 300 LLA LHI 003 LAM NDA JFZ SAVPOS JMP NOGDY
003 003 003	102 103 104 106 110 111 114 117	$\begin{array}{c} 041 \\ 315 \end{array}$		001 001	NOGDX, NOGDY,	INB INB LAB NDI 017 ORI 260 DCC JFZ TRYSEC LXH 201 001 CAL MSG JMP OVER
003 003	130 131 132 133 135 136 137	$\begin{array}{c} 041 \\ 126 \\ 054 \\ 136 \\ 056 \\ 162 \\ 054 \\ 163 \\ 311 \end{array}$	374 372	001	SAVPOS,	LXH 374 001 LDM INL LEM LLI 372 LMD INL LME RET
003 003 003 003 003 003 003	261 262 263 264 265	377 001 000 001 001 377				377 001 000 001 001 001 377

003	267	000	000
003	270	001	001
003	271	000	000
003	272	377	377
003	273	377	377
003	274	000	000
003	275	377	377
003	276	001	001
003	277	377	377

004 000

CKINP,

004 020

INPUTN,

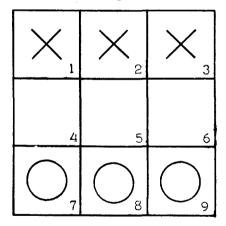
004 200

PRINT,

HEXPAWN - A MINI CHESS GAME

The possibility of playing a game of chess against a computer has undoubtably crossed the minds of most people that have had exposure to computers in one way or another. However, the game's near-limitless number of possible board configurations and moves makes it impossible to program on a small computer system. An alternative is to simplify the game to allow it to be programmed for the small computer system. Hexpawn is one such game.

Hexpawn consists of a 3×3 playing board and six pawns, three pawns for each player. The starting configuration is illustrated below. The pawns move in a manner similar to their moves in chess. A pawn can move one square forward, provided the square it is moving to is vacant, or one square diagonally to capture an opponent's pawn. A diagonal move cannot be made if an opponent's pawn is not captured by the move. The object of the game is to move a pawn to the opponent's side of the board while blocking the opponent from doing so, or capture all of the opponent's pawns. A game is a draw when no one can make a legal move.



The game starts by the current board configuration being printed followed by a request for a human to enter the first move. Each move is made by entering the number of the square which contains the pawn to be moved, followed by the number of the square to which the pawn is to be moved. The computer then makes its move, and prints the new board configuration. It then waits for the human's next move. After each move by the human and by the computer, the board is examined to determine if the game has been won by either side. When this occurs, an appropriate message is printed to indicate the end of the current game, and a new game is started. Should the human make an illegal move, the computer rejects the move and requests a new one.

As one can see, the game is fairly simple and requires no more than three or four moves by each side to complete. Thus, to make the game interesting, the program is written to provide the computer with ARTIFICIAL INTELLIGENCE. The computer is given the ability to decide which move it should make in an effort to win the game. That is, after each move is made by the human player, the computer examines the board and decides which, of all possible moves, it will make. If a move is made by the computer which results in the computer losing the game, that move is noted as an undesired move which should not be made again when that same board configuration if encountered. Thus, the computer learns from its mistakes, and eventually becomes so efficient in its ability to play the game that the best one can hope for is to play to a draw with the computer.

This version of Hexpawn is written to reside in five pages (256 bytes per page) of an 8008 or 8080 microcomputer system. The program may be reduced somewhat be revising or removing some of the text messages, if the user is limited in the amount of memory available. There are also several portions of the program which could be rearranged into subroutines, allowing for further compression of the program. Also, the table which is used to restore the program to its initial state of ignorance may be deleted along with the associated restoration program steps. If this is done, the program can be restored by simply reloading the program into memory. Making such changes to the program can reduce the memory required to less than four pages. The program was written to make it easy to follow the logic rather than minimize the amount of memory required. The reader can see that the memory required is, however, considerably less than that needed to duplicate the same game using a higher level language.

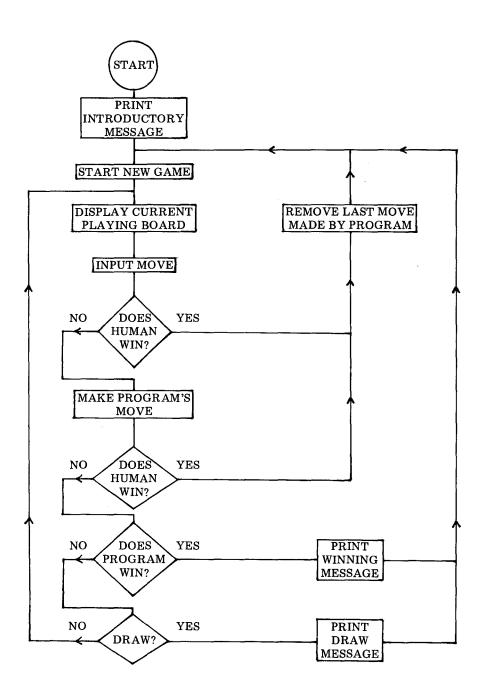
The flow chart on the following page illustrates the basic flow of the program. As one may observe, the game progresses in the same manner as a chess game. Each side makes one move at a time, and checks the board at the completion of each move to determine whether there is a winner. A verbal description of the flow chart will now be presented.

FUNDAMENTAL OPERATION OF THE HEXPAWN PROGRAM

The program starts by printing an introductory message which describes the operation of the game for a person that is playing the game for the first time. A game is then started by displaying the playing board along with the opening positions of the pawns for the human player to examine.

The program next requests the human to input a move. When the move has been received, it is checked to determine whether the player's move was to the opposite side of the board. That would indicate the challenger had won the game. If not, the move is entered on the current board. The program then examines the board and determines which move it will make in response to the player's move. If the human's move was a winning move, the program will remove the last move that it made from its list of possible moves. Since the last move that the computer made resulted in a win by the human, it does not want to make the same mistake twice. One may note that it requires at least two moves by a human to win a game, so that there will always be an initial move by the program which can be deleted.

When the program examines the current playing board, it selects a move from a list of possible moves which it may make for the current board configuration. If this list has had all of its entries removed, because the human has won as a result of making those moves, the game is conceded to the human and the move that the program just made will be removed. Otherwise, the program makes the move indicated in the list. It then determines whether its move has won the game or has resulted in a draw (all remaining pawns are blocked from making a move). If the game is a win or draw, an appropriate message is printed. A new game is then started. If not, the game is continued and the program returns to print the current



2 - 4

playing board for the player to examine.

TABLES USED BY THE HEXPAWN PROGRAM

The most important portion of a program such as Hexpawn is that which decides what move is to be made for a specific board configuration. This operation is performed through the use of four tables in this program. These tables are used to: Find the matching board configuration, direct the program to the list of possible moves for each configuration, select the move to be made, and provide the actual codes for making the move. Each of these tables are presented in the listing at the end of this chapter. Due to the size of these tables, only sample entries of each will be presented in the following discussion.

The MODEL table is a table used by the program to determine the current board configuration. It consists of 33 pairs of bytes which define all the possible board configurations immediately following a human's move. The first byte of each pair indicates the positions of the program's pawns on the board. For each pawn in a square, the bit corresponding to that square is set to '1.' If a pawn is not in a square, the bit corresponding to that square is '0.' The squares of the board defined in the first byte are as follows:

BIT POSITION	B 7	B6	B5	B 4	B3	B2	B1	B0
BOARD POSITION	1	2	3	4	5	6	Х	Х

The reader may note that bits B1 and B0 do not have any position on the playing board defined for them. The reason for this is that if any of the program's pawns reach position 7, 8, or 9, the game will be over with the computer winning. It will not be necessary to store the fact that the program's pawn has reached the last row. Consequently, bits B1 and B0 will always be set to a zero. The same is true for the human's pawns, which are defined by the following bit definitions in the second byte of the model pair.

BIT POSITION	B7	B6	B5	B 4	B3	B2	B1	B0
BOARD POSITION	Х	Х	4	5	6	7	8	9

2 - 5

Bits B7 and B6 are not defined for any position on the board, since a move to position 1, 2, or 3 is a winning move for the human. Bits B7 and B6 of the second byte are always a zero.

For example, suppose the first move made by the human was from square 7 to square 4. The current board configuration would be represented by the following byte pair (using octal notation):

PGM'S PAWNS	340
HUMAN'S PAWNS	043

This corresponds to a physical board configuration of:

X | X | X 0| | |0|0

When the program finds the byte pair in the model table that matches the current board, it sets up a pointer to the MODEL-TO-MOVE INDEX table. The model-to-move index table consists of a list of pointers. Each entry in this table points of a list of possible moves for the current board configuration. The list of moves is contained in the MOVE INDEX table.

The move index table contains a list of moves which may be made for each of the 33 models in the model table. Each list contains from one to three numbers which indicate a possible move. Each list is terminated by a 200 octal entry. The move numbers range in value from 1 to 15, and indicate possible moves taken from the list on the following page.

The move number is a number which is contained in the move index table. The number in the FROM column is the square on the playing board that the program's pawn is to be moved from. The TO column contains the square that the program's pawn is to be moved to. The result of the move is indicated in the last column. When a number of a move is read from the move index table, it is used to set up a pointer to the MOVE table.

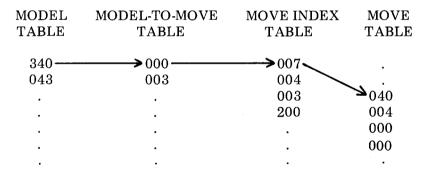
The move table contains a four byte grouping for each of the

MOVE NO.	FROM	ТО	RESULT
1	1	4	
2	1	5	CAPTURE
3	2	4	CAPTURE
4	2	5	••-
5	2	6	CAPTURE
6	3	5	CAPTURE
7	3	6	
8	4	7	COMPUTER WINS
9	4	8	COMPUTER WINS
10	5	7	COMPUTER WINS
11	5	8	COMPUTER WINS
12	5	9	COMPUTER WINS
13	6	8	COMPUTER WINS
14	6	9	COMPUTER WINS
15	-	-	DRAW

moves in the table above. The first byte has one bit set. This bit represents the original position of the program's pawn which is to be moved (as discussed previously). The second byte has a bit set to represent the position to which the program's pawn will be moved. The third byte indicates the position of the player's pawn which will be captured by the move, if the move results in a capture. If no capture will be made, the third byte will be all zeros. The fourth byte depicts the result of the move. If the fourth byte is zero, the program will continue with the game by requesting a move by the human. If the move results in the computer winning, the fourth byte will contain the ASCII code for the number of the square the computer will move into to win the game. This code is used in setting up the winning message. If the move will result in a draw, the fourth byte will contain the octal value 100. The following example contains the four byte group used to define move number three. The reader may note that bytes 2 and 3 both correspond to square 4 of the playing board. Byte 2 indicates the program's pawn moving to square 4, and byte 3 indicates the human's pawn at square 4 being captured by the move.

MOVE NO. 3	BYTE NO. 1	100
	BYTE NO. 2	020
	BYTE NO. 3	040
	BYTE NO. 4	000

The move index table is the means by which the program learns to play the game. Each time a move is selected from the move index table that location in the table is saved. After every move, a test for a possible win by the human is made. If the human wins, the location in the move index table which was saved by the program is zeroed. Once zeroed, the program will not be able to make the same move. The program will skip that location and make the next move indicated in the list. For example, suppose in the first game the player's first move was from 7 to 4. The program would find the first model in the model table as a match. It would fetch the pointer from the first location in the model-to-move index table and select move 7 as its first move. This sequence is illustrated below.



From this sequence it may be observed that the program has moved from square 3 to 6. This move allows the human to win easily by moving from square 4 to 2 thus capturing the program's pawn. The program then removes 007 from the list of moves for the first model by loading a 000 in the first location of that list. Now, if the human makes the same opening move of square 7 to square 4, the program will skip the first location in the list of moves for the first model, since it is zero. It will then make the second move, which is move number 4. Thus, the program has learned that move number 7 was not the proper move to make for that particular board model.

Another manner in which the program learns is when all the moves for a specific model have been zeroed. When the program searches the move index table for a move, and it reaches the 200 byte, which terminates the list, it concedes the game to the human. The program knows that on the next move, no matter what move it makes, it will lose the game. At this point the program will go back to the previous move that it made before getting into the predicament and zero it in the move index table. In this way, the program is prevented from making the same move which brought it to the point of having to concede the game.

TEXT MESSAGES USED IN HEXPAWN

There are several messages used in Hexpawn to inform the human player of the setup of the game, to indicate when the human has input an illegal move, to display the current board configuration, and to signify the outcome of the game. These messages are of variable length, and may require more than one line of output. The content of these messages may be altered by the reader, if desired, to reflect greater emotion by the program at winning and losing. The messages presented here were kept fairly low key to conserve memory space, since the operating program and associated tables alone require almost 1K bytes of memory. The message strings are presented next.

> "HERE'S THE BOARD 1|2|3 4|5|6 7|8|9 I'M X, YOU'RE O MOVE ONE SQUARE FORWARD IF VACANT OR ONE SQUARE DIAGONALLY TO CAPTURE YOU START."

"NO ! NO !"

"I CONCEDE!"

"YOU WIN."

"NO GOOD! MUST START AGAIN"

"DRAW, NO ONE WINS."

"I MOVE TO . I WIN! YOU LOSE!"

"I WIN! YOU HAVE NO PAWNS."

These text messages are stored as a continuous string of ASCII characters with each message terminated by a zero byte. The playing board, however, is stored on the same page as the model table and temporary data, as it is updated after each move by the program. The board output, FROM and TO messages are stored as shown next.

"X ! X ! X ! ! O ! O ! O FM " TO "

The program uses a common subroutine to output these messages to the user's output device. This subroutine is called with memory pointer registers H and L set to the starting address of a message. It fetches each character from the storage area and presents it to the user's output routine in the accumulator. When it encounters the zero byte, it returns to the calling program. The listing for this routine is presented next. It is labeled MSG.

MSG,	LAM	Fetch character to print
	NDA	End of message?
	RTZ	Yes, return
	CAL PRINT	No, print character
1 C	INL	Incr low addr msg pointer
	JFZ MSG	If non-0, continue output
	INH	Else, incr page addr pointer
	JMP MSG	

As one can see, the MSG subroutine is quite straight-forward. It simply fetches characters from memory starting at the location set up by the calling program in the H and L registers, and calls the user provided subroutine PRINT to output the character to the system output device. The PRINT subroutine must take the character in the accumulator and perform whatever is required to output that character to the printer or display device. This routine is free to use registers B through E in outputting the character. The only requirement is that if registers H and L must be used, they must be restored to their initial value before returning to the MSG subroutine.

The MSG subroutine is called to output every message by this program. Therefore, in order to change the messages, the reader simply stores the ASCII codes for the messages desired in a continuous string in memory, and then stores a zero byte to terminate the message. To output a message, the program sets the pointer to the starting address of the message and calls the MSG subroutine.

THE HEXPAWN PROGRAM

The reader may refer to the flow chart presented previously during the following discussion of the operating portion of the program.

The first part of the Hexpawn program outputs the introductory message and resets the move index table to its initial state of intelligence. The move index table is initialized by transferring the contents of the RESTORE MOVE list into the ACTIVE MOVE list. The restore move list is a copy of the move index table with all the possible moves contained in it. The active move list is actually the move index table, which will have its contents zeroed as the program learns to play the game. This restore routine transfers the upper half of page 04, which contains the restore list and the beginning of the text messages, down to the lower half of page 04, which contains the active move list, or move index table. The actual list, however, only requires approximately 3/8 of the page. This should be noted so that one does not try to store messages in the area from 137 through 177 on page 04.

This restore routine is performed only when the program is started at the beginning. After each game is completed, the program returns to the next section, starting at the label AGAIN so that the move index table will not be reset. The only way that the move index table can be reset is for the operator to restart the program.

The listing for the initial portion of the Hexpawn program is presented below.

START,	LLI 076 LHI 005	Set pointer to intro. msg
	CAL MSG	Print introduction
	LHI 004	Set pointer to move index pg
	LDI 000	Init. actv move list pointer
	LEI 200	Init. rstr move list pointer
RSTR,	LLE	Set restore list pointer
	LAM	Fetch restore list entry
	LLD	Set pointer to active list
	LMA	Store entry in active list
	IND	Increment active list pointer
	INE	Increment restore list pointer
	JFZ RSTR	Done? No, cont. transfer

The next section of the program is the one which prepares the board output to display the current setup of the playing board. There are two points at which the program enters this routine. The first is at the instruction labeled AGAIN. This entry point resets the playing board, as stored in locations 03 000 for the X pawns, and 03 001 for the O pawns. This is the starting setup as shown in the figure on the first page of this chapter. It then proceeds to the other entry point of this routine. The second entry point is labeled PBD. This point is entered when the program is in the middle of a game. This portion of the routine sets up the board output message to display the current positions of the pawns in the following manner.

First, the board output message is cleared by storing space characters, ASCII code 240, at the locations in the board output that represent the possible pawn positions. The current positions of the X pawns and O pawns are then determined by the subroutines STX and STO.

When the STX subroutine is first called, the accumulator contains

the current X board as stored in location 03 000. This is rotated left one bit to load the CARRY with the condition of square 1 with respect to the presence of an X pawn. If the carry is set to 1, the ASCII code for an X is stored in the location of square 1 in the board output message. If the carry is reset to 0, the routine simply returns and the contents of square 1 remain a space character. The next time STX is called, the accumulator contains the current X board rotated to the left one bit so that when it is rotated to the left again the carry will indicate the presence, or absence, of an X pawn for square 2. Each time the STX routine is called, the memory pointer registers H and L are set to indicate the location in the board output message where the X is to be stored if it is present at that location. The STO subroutine stores the ASCII code for the 0 character in the locations in the board output message where there are 0 pawns present in a manner similar to the STX subroutine. When the board output message is set to reflect the current board set up, the MSG subroutine is called to display the board for the player to examine. The following is the listing of this board set up and display routine.

AGAIN,	LLI 000 LHI 003	Set pointer to current board
	LMI 340 INL	Set board to starting setup
	LMI 007	
PBD,	LLI 302	Set pointer to brd printout
	LBI 240	Set space char to clear board
	LMB	Store space in '1'
	LLI 304	
	LMB	Store space in '2'
	LLI 306	
	LMB	Store space in '3'
	LLI 311	
	LMB	Store space in '4'
	LLI 313	
	LMB	Store space in '5'
	LLI 315	
	\mathbf{LMB}	Store space in '6'
	LLI 320	
	LMB	Store space in '7'
	LLI 322	

LMB	Store space in '8'
LLI 324	-
LMB	Store space in '9'
LLI 000	Set pointer to current X board
LAM	Fetch X board
INL	Advance to O board
LBM	Fetch current O board
LLI 302	Set pointer to 1 position
CAL STX	If X here, store character
LLI 304	Set pointer to 2 position
CAL STX	If X here, store character
LLI 306	Set pointer to 3 position
CAL STX	If X here, store character
LLI 311	Set pointer to 4 position
CAL STX	If X here, store character
LCA	Save X board
LAB	Fetch O board
RLC	Position to 4
RLC	
CAL STO	If O here, store character
LBA	Save O board
LLI 313	Set pointer to 5 position
LAC	Fetch X board
CAL STX	If X here, store character
LCA	Save X board
LAB	Fetch O board
CAL STO	If O here, store character
LBA	Save O board
LLI 315	Set pointer to 6 position
LAC	Fetch X board
CAL STX	If X here, store character
LAB	Fetch O board
CAL STO	If O here, store character
LLI 320	Set pointer to 7 position
CAL STO	If O here, store character
LLI 322	Set pointer to 8 position
CAL STO	If O here, store character
LLI 324	Set pointer to 9 position
CAL STO	If O here, store character
LLI 300	-
11 200	Set pointer to board printout

	CAL MSG	Print current board
STX,	RLC RFC LMI 330 RET	Bit set? No, return Yes, put X in board
STO,	RLC RFC LMI 317 RET	Bit set? No, return Yes, put O in board

After the current board is outputted, the program requests the player to enter a move by first entering the number of the square which contains the pawn to be moved, and then the number of the square to which the pawn is to be moved. The input request is indicated by the output of the message FM which is output as part of the board output message. The program then calls the user supplied input routine to accept a character from the system input device.

The input subroutine is a user provided routine which must input a character from the keyboard device and return with the ASCII code for that character in the accumulator. This subroutine is free to use registers A through E in inputting the character. If registers H and L are required to be used, they must be restored to their original contents before returning to the calling program. If the system's input device does not provide automatic echo of the inputted character to the output device, this input routine should include some provision for echoing the character received to the output device. This subroutine is called only in the move input routine being presented here.

When the FROM square is received, it is checked first to determine whether it is a valid number from 1 to 9, since these are the only valid entries expected. This is checked by calling the FNUM subroutine. If the input is not within these limits, the ERROR routine is entered. The error routine is called at several points in the next group of routines whenever the move which has been input is found to be illegal. The error routine prints the message "NO! NO!" and then jumps to the PBD entry point of the program to request a new input.

If the FROM input is valid, the 260 portion of the ASCII code is removed, and the binary value of the number entered is stored in location 03 002. The current O board is then checked to determine whether an O pawn does reside in the square designated by the input. The binary value of the FROM square is used as a counter by the RTAL subroutine which rotates the current O board until the carry bit indicates the presence or absence of an O pawn in that position. If there is no O pawn, the error routine is entered.

The message TO is then output and the INPUT routine is called to input the square to which the pawn is to be moved. This input is also checked by calling the FNUM routine to determine whether it is within the limits of the expected input. If it is valid, it is changed to its binary value and stored in location 03 003. The program then proceeds to the next routine which checks that the move is legal.

The listing for this portion of the program is presented below.

Set pointer to input storage
Input FM move
Save input
Number valid?
No, error
Fetch number
Delete ASCII code
Save FM location
Save bit count for RTAL
Set pointer to O board
Fetch O board
Is pawn in FM position?
No, illegal move
Set pointer to TO msg
Print TO
Set pointer to input storage
Input TO move
Save TO input

	CAL FNUM JTS ERROR LAM NDI 017 LMA	Input valid? No, error Fetch number Delete ASCII code Save TO location
ERROR,	LLI 364 LHI 005	Set pointer to error message
	CAL MSG LHI 003	Print error message
	JMP PBD	Print current board
RTAL,	DCB	Decrement bit count
	RTZ	If zero, return
	RLC	Else, rotate left
	JMP RTAL	
FNUM,	LAM	Fetch ASCII number
	CPI 261	Is number valid?
	RTS	No, return with S flag set
	SUI 272	If number is valid, return
	ADI 200 RET	With S flag set

Once the FROM and TO values are received, the move must be checked to determine whether it is legal. First, a move forward is checked by subtracting 3 from the FROM value and checking it with the value stored for the TO value. If the move is forward one square, the position of the X pawns must be checked to make sure an X pawn is not blocking the move. The BLK routine is entered to check the forward move. BLK sets the TO value as a counter and fetches the current X board, which is then rotated left by the RTAL subroutine, placing the bit corresponding to the location the O pawn is to be moved, to the sign position in the accumulator. If this bit is set, the move is blocked by an X pawn and the error routine is entered. If not, the BLK routine returns to the mainstream of the program at the HMV label to make the move as entered.

If the move is not forward, a diagonal move to the left or right is

examined. By adding 1 to the forward move, the new value indicates a diagonal move to the right. If this matches the TO value, and is not equal to 7, which would be an illegal move from 9 to 7, the capture of an X pawn is checked, since a diagonal move must capture an opponent's pawn. If the move is not to the right, 2 is subtracted from the forward move and a diagonal move to the left is tested. If this matches and it is not equal to 3, indicating an illegal move from 7 to 3, an X pawn capture will be checked. If any of the above illegal conditions occur, the error routine is entered. When a capture move is indicated, the bit position in the current X board of the bit to be deleted is set up by the RTLP subroutine. The presence of an X pawn is checked, and if not there, the error routine is entered.

Once the preliminaries are complete, the move is checked for a win by the human player. If the move is to square 1, 2, or 3, the human has won the game. The HWIN routine is then entered to perform the required steps to teach the program. This is accomplished by zeroing the last move made by the program in the move index table. The HWIN routine then outputs a congratulatory message and starts a new game.

If the move does not result in a win, the move is entered in the current O board by resetting the bit indicating the FROM position and setting the bit indicating the TO position. If a capture was made, the bit in the current X board in the TO position is reset.

The listing for this routine is shown below.

DCL	Set FM pointer
LAM	Fetch FM
SUI 003	Is move forward?
INL	Check against TO
CPM	
JTZ BLK	Yes, check if legal
ADI 001	No, move right 1 square
CPM	Is TO here
JTZ CKCAP	Yes, check for capture
SUI 002	No, move left 1 square
CPI 003	Is move from 7 to 3?

CKCAP,	JTZ ERROR CPM JFZ ERROR CPI 007 JTZ ERROR LAI 200 CAL RTLP LEA	Yes, illegal Is TO here? No, illegal move Is move to 7? Yes, error Fetch TO move Set up to calculate capture Bit by rotating right Save capture bit
	LLI 000 NDM	Set X board pointer Capture?
	JTZ ERROR	No, illegal move
HMV,	LLI 002	Set pointer to FM
	LAM	Fetch FM location
	CAL RTAR	Set up FM bit
	LDA	Save FM bit
	INL	Set pointer to TO
	LAM	Fetch TO location
	CPI 004	Human wins?
	JTC HWIN	Yes, zero last move
	CAL RTAR	Set up TO bit Save TO bit
	LCA LLI 001	
	LAM	Set pointer to current O board Fetch current board
	XRD	Clear old set
	ORC	Set new position
	LCA	Save new O board
	LMA	Save current board
	DCL	Save carrent scara
	LAE	Fetch capture bit
	NDA	Capture?
	JTZ NOCP	No, skip
	XRM	Yes, delete piece
	LMA	Save current X board
RTAR,	LBA	Set bit count
	LAI 001	Set bit to rotate
RTLP,	DCB	Decrement bit count
	RTZ	If zero, return
	RRC	Else, rotate right

JMP RTLP

BLK,	LBM LLI 000 LAM CAL RTAL	Fetch TO move Set pointer to X board Fetch X board Check for blocked move
SET,	NDA	Is move blocked?
	JTS ERROR	Yes, illegal move
	LEI 000	Set for no capture
	JMP HMV	Return to make human move
HWIN,	LLI 004 LHI 003	Set pointer to last move
	LLM LHI 004	Fetch last move address
	LMI 000	Zero last move
	LLI 315 LHI 005	Set pointer to lose message
	CAL MSG JMP AGAIN	Print lose message Start new game
	JULI AGAIN	Start new game

Now that the human's move has been entered, it is time for the program to show what it knows about the game. This is the portion of the program which performs the table search and makes the resultant move that the program believes correct at the time for the given board configuration. The program first searches the model table for a matching model. If none is found, it is assumed that the move just input by the human is invalid. Some of the conditions which were not checked in the move input routine include a move forward into a square already occupied by an O pawn. (This results in the human eliminating one of his own pawns.) Or, a move from square 6 to square 4 in which an X pawn is captured. These moves will slip through the initial validity tests, but they do result in illegal board setups which are caught here. The only recourse for the program at this time is to start the game over again because the current board is unrecognizible.

When a model is found where the current X board and current O board match a byte pair in the model table, a pointer is calculated

from the relative position in the model table to the model-to-move index table. This pointer is calculated by dividing the low address of the X byte of the matching model by 2 and adding 106 to the result. The pointer is used to fetch the starting address of the list of possible moves in the move index table from the model-to-move index table.

The designated list of moves is then examined, and the first nonzero entry is used as the move the program will make in response to the current board model. If the program encounters a 200 byte in searching the table, it jumps to the ONO (OH! NO!) routine. Reaching a 200 byte indicates the program has made every move it can for the model, and they all lead to defeat. The ONO routine concedes defeat and also goes to the HWIN routine to eliminate the previous move with the intent that this model will not be encountered again.

When a move is found, a pointer is set up using the move number and the four bytes of the move are fetched from the move table. The last byte of this group is examined first to determine whether the program has won the game, indicated by the sign bit set in the last byte. Or, if the game is a draw, indicated by a 100 stored in the last byte. If the game is won by the computer, the WIN routine is entered to print the winning message and start a new game. If the game is a draw, the DRAW routine is entered, to print the draw message and start a new game. If the last byte is zero, the move is made as indicated by the first three bytes. The first byte has the bit set which is the FROM location of the X move, the second byte has the bit set which is the TO location of the X board and the third byte has the bit set which indicates which location in the O board has been captured. If the third byte is zero, the move does not capture any O pawn, and the game is continued by jumping to PBD.

When a capture is made, the O board is checked to determine whether all the O pawns have been captured. If so, the program has won the game. At this point, the program deletes the move that it has just made from the move table because it knows that if there was only one O pawn on the board, the program has a move open which will allow it to win by moving to the opponent's side rather than gobbling up the last pawn. The program will then print a message to inform the human that he has no more pawns! The listing for this final routine of the Hexpawn program is presented next. The reader will note the common MSG call followed by a jump to AGAIN which starts a new game. This instruction pair, labeled CMSG, was set up to conserve program space, as it is a common sequence used by several routines to print game concluding messages and then begin a new game.

NOCP,	LDM	Save new X board
CMDI	LLI 010	Set pointer to model table Fetch X board
SMDL,	LAD	
	CPM KPZ OLL F	X board match model?
	JTZ OHLF	Yes, try O half
CMD1	INL	Advance table pointer
SMD1,	INL	
	LAL	Check for end of table
	CPI 112	End of table?
	JFZ SMDL	No, continue search
	LLI 340	No match, illegal move made
~ ~ ~ ~	LHI 004	Print "NO GOOD!"
CMSG,	CAL MSG	Print message
	JMP AGAIN	Start new game
OHLF,	INL	Advance pointer to O board
ondr,		Fetch current O board
	CPM	O boards match?
	JFZ SMD1	
	JFZ SMDI	No, continue search
	DCL	Move pointer to X board
	LAL	Set up to calculate pointer
	RRC	Divide by 2
	ADI 106	Add to start of mdl index tbl
ŵ	LLA	Set pointer to mdl index tbl
	LLM	Fetch pntr to move index tbl
	LHI 004	Set pntr to move index table
MFD1,	LAM	Fetch move number
·	NDA	Move number here?
	JTS ONO	No move avail. Human wins
	JFZ MOVE	Move found, make it
	INL	Move zeroed, try next location

MOVE,	LEL	Save move location
	LLI 004	Set pointer to last move
	LHI 003	-
	LME	Save location as last move
	RLC	Set up pointer to move
	\mathbf{RLC}	Storage table
	ADI 174	
	LLA	Set pointer
	LDM	Fetch FM bit
	INL	Advance pointer
	LCM	Fetch TO bit
	INL	Advance pointer
	LEM	Fetch capture bit
	INL	Advance pointer
	LAM	Fetch contest bit
	NDA	Is game over?
	JTS WIN	Yes, computer wins
	JFZ DRAW	Yes, draw
	LLI 000	Set pointer to X board
	LAM	Fetch current X board
	XRD	Clear old position
	ORC	Set new position
	LMA	Save new X board
	INL	Advance pointer to O board
	LAE	Fetch capture bit
	NDA	Capture?
	JTZ PBD	No, print new board
	XRM	Yes, delete piece
	LMA	Save new O board
	JFZ PBD	Non-0, continue game
	LLI 004	Set pointer to last move
	LLM	Fetch last move location
	LHI 004	Set pointer to active move list
	LMI 000	Cancel last move
	LHI 005	Set pointer to msg ''I WIN
	LLI 330	YOU HAVE NO PAWNS!"
	CAL CMSG	Print msg and start again
		0

WIN,	LLI 025 LHI 005	Set pointer to store win move			
	LMA	Store win move in message			
	LLI 011	Print "I MOVE TO ."			
	JMP CMSG	"I WIN, YOU LOSE"			
DRAW,	LLI 052 LHI 005	Prnt 'DRAW, NO ONE WINS'			
	JMP CMSG	Print draw message			
ONO,	LLI 374 LHI 004	Print "I CONCEDE!"			
	CAL MSG	Then zero last move			

Well! That's it! Now the Hexpawn program is presented in its final assembled form to be loaded into an 8008 based microcomputer system. The operating portion of the program resides on pages 01 and 02 and the tables and messages are on pages 03 through 05. Due to the length of assembled listings for the tables and messages, they will be presented as an octal dump.

001	000	066	076		START,	LLI 076
001	002	056	005			LHI 005
001	004	106	171	002		CAL MSG
001	007	056	004			LHI 004
001	011	036	000			LDI 000
001	013	046	200			LEI 200
001	015	364			RSTR,	LLE 🛸
001	016	307				LAM
001	017	363				LLD
001	020	370				LMA
001	021	030				IND
001	022	040				INE
001	023	110	015	001		JFZ RSTR
001	026	066	000		AGAIN,	LLI 000
001	030	056	003			LHI 003
001	032	076	340			LMI 340

	034	060				INL
	035	076				LMI 007
	037	066	302		PBD,	LLI 302
001	041	016	240			LBI 240
001	043	371				LMB
001	044	066	304			LLI 304
001	046	371				LMB
001	047	066	306			LLI 306
001	051	371				LMB
001	052	066	311			LLI 311
001	054	371				LMB
001	055	066	313			LLI 313
001	057	371				LMB
001	060	066	315			LLI 315
001	062	371				LMB
001	063	066	320			LLI 320
001	065	371				LMB
001	066	066	322			LLI 322
001	070	371				LMB
	071	066	324			LLI 324
001	073	371				LMB
001	074	066	000			LLI 000
001	076	307				LAM
	077	060				INL
001	100	317				LBM
001	101	066	302			LLI 302
001	103	106	157	002		CAL STX
001	106	066	304			LLI 304
001	110	106	157	002		CAL STX
001	113	066	306			LLI 306
001	115		157	002		CAL STX
001	120	066	311			LLI 311
001	122	106	157	002		CAL STX
001	125	320				LCA
001	126	301				LAB
001	127	002				RLC
001	130	002				RLC
001	131	106	164	002		CAL STO
001	134	310				LBA
001	135	066	313			LLI 313

001	137	302				
001	140	106	157	002		
001	143	320				
001	144	301				
001	145	106	164	002		
001	150	310				
001	151	066	315			
	153	302				
001	154	106	157	002		
001	157	301				
001	160	106	164	002		
001	163	066	320			
	165	106	164	002		
	170	066	322			
001	172	106	164	002		
001	175	066	324			
$\begin{array}{c} 001 \\ 001 \end{array}$	177	106	164	002		
001	202	066	300			
001	204	106	171	002		
001	207	066	002			
	211	106		006		
001		370				
001	215		303	002		
001	220		207			
	223	307				
	224	044	017			
	226	370				
001	227	310				
	230	061				
001	231	307				
	232	106	223	002		
	235		207	002		
	240		333			
001	242	106	171	002		
	245		003			
	247	106	000	006		
	252	370				
	253		303	002		
001	256	160		002		
001		307				
		-				

LAC CAL STX LCA LAB CAL STO LBA LLI 315 LAC CAL STX LAB CAL STO LLI 320 CAL STO LLI 322 CAL STO LLI 324 CAL STO LLI 300 CAL MSG LLI 002 CAL INPUT LMA CAL FNUM JTS ERROR LAM NDI 017 LMA LBA DCL LAM CAL RTAL JFC ERROR LLI 333 CAL MSG LLI 003 CAL INPUT LMA CAL FNUM JTS ERROR LAM

001	l 262	044 01	17		NDI 017
001	26 4	370			LMA
001	l 265	061			DCL
001	266	307			LAM
001	267	024 00)3		SUI 003
001	l 271	060			INL
001	L 272	277			CPM
001	1 273	$150 2^{2}$	42 002		JTZ BLK
001	L 276	004 00)1		ADI 001
	L 300	277			CPM
	l 301		17 001		JTZ CKCAP
	l 304	024 00			SUI 002
	1 306	074 00			CPI 003
	1 310		07 002		JTZ ERROR
	313	277			CPM
	l 314	110 20	07 002		JFZ ERROR
00	1 317	074 00	07	CKCAP,	CPI 007
001	l 321	150 20	07 002	,	JTZ ERROR
001	l 324	317			LBM
001	325	006 20	00		LAI 200
001	l 327	106 23	34 002		CAL RTLP
001	l 332	340			LEA
001	L 333	066 00	00		LLI 000
001	1 335	247			NDM
001	L 336	150 20	07 002		JTZ ERROR
001	l 341	066 00)2	HMV,	LLI 002
001	l 343	307			LAM
001	l 344	106 23	81 002		CAL RTAR
001	l 347	330			LDA
001	L 350	060			INL
001	l 351	307			LAM
001	l 352	074 00)4		CPI 004
001	l 354	140 32	23 002		JTC HWIN
001	1 357	106 23	31 002		CAL RTAR
001	l 362	320			LCA
001	L 363	066 00	01		LLI 001
001	L 365	307			LAM
001	L 366	253			XRD
001	l 367	262			ORC
001	L 370	320			LCA

001	371	370					LMA
001	372	061					DCL
001	373	304					LAE
001	374	240					NDA
001	375	150	002	002			JTZ NOCP
002	000	257					XRM
002	001	370					LMA
002	002	337			Ν	NOCP,	LDM
002	003	066	010				LLI 010
002	005	303			S	MDL,	LAD
002	006	277					CPM
002	007	150	034	002			JTZ OHLF
002	012	060					INL
002	013	060			S	SMD1,	INL
002	014	306					LAL
002	015	074	112				CPI 112
002	017	110	005	002			JFZ SMDL
002	022	066	340				LLI 340
002		056	004				LHI 004
002	026	106	171	002	C	CMSG,	CAL MSG
002	031	104	026	001			JMP AGAIN
002	034						
002	034	060			C)HLF,	INL
002	035	302					LAC
002	036	277					CPM
002	037	110	013	002			JFZ SMD1
002	042						
002	042	061					DCL
002	043	306					LAL
002	044	012					RRC
002	045	004	106				ADI 106
002	047	360					LLA
002	050	367					LLM
002	051	056	004				LHI 004
002	053	307			Ν	AFD1,	LAM
002	054	240					NDA
002	055	160	314	002			JTS ONO
002	060	110	067	002			JFZ MOVE
002	063	060					INL
002	064	104	053	002			JMP MFD1

	067	346			MOVE,	\mathbf{LEL}
002	070	066	004			LLI 004
002	072	056	003			LHI 003
002	074	374				LME
002	075	002				RLC
002	076	002				RLC
002	077	004	174			ADI 174
002	101	360				LLA
002	102	337				LDM
002	103	060				INL
002	104	327				LCM
002	105	060				INL
002	106	347				LEM
002	107	060				INL
002	110	307				LAM
002	111	240				NDA
002	112	160	262	002		JTS WIN
002	115	110	274	002		JFZ DRAW
002	120	066	000			LLI 000
002	122	307				LAM
002	123	253				XRD
002	124	262				ORC
002	125	370				LMA
002	126	060				INL
002	127	304				LAE
002	130	240				NDA
002	131	150	037	001		JTZ PBD
002	134	257				XRM
002	135	370				LMA
002	136	110	037	001		JFZ PBD
002	141					
002	141	066	004			LLI 004
002	143	367				LLM
002	144		004			LHI 004
	146	076				LMI 000
	150	056				LHI 005
	152	066				LLI 330
	154		026	002		CAL CMSG
	157		•			
	157	002			STX,	RLC
0.04	<u> </u>				S121,	

002	160	003				RFC
002	161	076	330			LMI 330
002	163	007				RET
002	164					
002	164	002			STO,	RLC
002	165	003				RFC
002	166	076	317			LMI 317
002	170	007				RET
002	171					
002	171	307			MSG,	LAM
002	172	240				NDA
002	173	053				RTZ
002	174	106	100	006		CAL PRINT
	177	060				INL
	200	110	171	002		JFZ MSG
	203	050				INH
002	204	104	171	002		JMP MSG
002	207					
002	207	066	364		ERROR,	LLI 364
002	211	056	005		,	LHI 005
002	213	106	171	002		CAL MSG
002	216	056	003			LHI 003
002	220	104	037	001		JMP PBD
002	223					
002	223	011			RTAL,	DCB
002	224	053			,	RTZ
002	225	002				RLC
002	226	104	223	002		JMP RTAL
002	231					
002	231	310			RTAR,	LBA
002	232	006	001		,	LAI 001
002	234	011			RTLP,	
002	235	053			,	RTZ
002	236	012				RRC
	237		234	002		JMP RTLP
	242			_		
	242	317			BLK,	LBM
	243		000			LLI 000
	245	307				LAM
	246		223	002		CAL RTAL
		100		004		

002		240			SET,	NDA
002			207	002		JTS ERROR
002		046				LEI 000
002		104	341	001		JMP HMV
002						
002		066	025		WIN,	LLI 025
002	264	056	005			LHI 005
002	266	370				LMA
002	267	066	011			LLI 011
002	271	104	026	002		JMP CMSG
002	274					
002	274	066	052		DRAW,	LLI 052
002	276	056	005			LHI 005
002	300	104	026	002		JMP CMSG
002	303					
002	303	307			FNUM,	LAM
002	304	074	261			CPI 261
002	306	063				RTS
002	307	024	272			SUI 272
002	311	004	200			ADI 200
002	313	007				RET
002	314					
002	314	066	374		ONO,	LLI 374
002	316	056	004		,	LHI 004
002	320	106	171	002		CAL MSG
003	323					
002	323	066	004		HWIN,	LLI 004
002	325	056	003			LHI 003
002	327	367				LLM
002	330	056	004			LHI 004
002	332	076	000			LMI 000
002	334	066	315			LLI 315
002	336	056	005			LHI 005
002	340	106	171	002		CAL MSG
002	343	104	026	001		JMP AGAIN
002	346					

TEMPORARY DATA

003 000

000 000 000 000 000

003 010	340	043	340	016	340	025	260	021
003 020	150	041	240	062	300	051	150	014
003 030	160	034	26 0	012	304	061	140	054
$003 \ 040$	140	021	140	024	240	041	070	010
003 050	200	070	120	030	104	060	230	010
003 060	240	014	210	042	054	040	060	020
003 070	110	040	110	010	220	020	044	020
003 100	200	060	240	032	100	020	250	052
003 110	040	070						

MODEL-TO-MOVE INDEX TABLE

003	112			000	004	010	013	017	023
003	120	027	033	036	041	043	046	051	054
003	130	057	061	063	065	070	073	075	077
003	140	101	103	107	112	115	120	123	125
003	150	131	133	135					

MOVE TABLE

,

003	200	200	020	000	000	200	010	020	000
003	210	100	020	040	000	100	010	000	000
003	220	100	004	010	000	040	010	020	000
003	230	040	004	000	000	020	000	000	267
003	240	020	000	000	270	010	000	000	267
003	250	010	000	000	270	010	000	000	271
003	260	004	000	000	270	004	000	000	271
003	270	000	000	000	100				

BOARD OUTPUT MESSAGE

003	300	215	212	330	336	330	336	330	215
003	310	212	240	336	240	336	240	215	212
003	320	317	336	317	336	317	215	212	306
003	330	315	240	000	240	324	317	240	000

MOVE INDEX TABLE

004 000	007	004	003	200	001	004	005	200
004 010	001	002	200	007	006	010	200	007
004 020	003	013	200	007	002	006	200	003
$004 \ 030$	004	017	200	005	012	200	005	006
$004 \ 040$	200	010	200	002	003	200	005	017
$004 \ 050$	200	006	017	200	006	007	200	017
$004 \ 060$	200	010	200	002	200	005	010	200
$004 \ 070$	003	016	200	013	200	017	200	017
004 100	200	016	200	007	006	010	200	003
$004 \ 110$	013	200	005	013	200	002	010	200
$004 \ 120$	006	016	200	002	200	001	006	017
004 130	200	017	200	017	200	006	200	

RESTORE LIST

004200004210004220004230004240004250004260004270	007 001 003 004 200 200 200 200 003	004 002 013 017 010 006 010 016	003 200 200 200 200 017 200 200	002 200	001 006 002 012 003 006 200 200	004 010 006 200 200 007 005 017	005 200 200 005 005 200 010 200	200 007 003 006 017 017 200 017
001 800	001			000			000	000
$004 \ 250$	200	006	017	200	006	007	200	017
$004 \ 260$	200	010	200	002	200	005	010	200
$004 \ 270$	003	016	200	013	200	017	200	017
004 300	200	016	200	007	006	010	200	003
$004 \ 310$	013	200	005	013	200	002	010	200
$004 \ 320$	006	016	200	002	200	001	006	017
004 330	200	017	200	017	200	006	200	

MESSAGE STORAGE

004	340	215	212	316	317	240	307	317	317
004	350	304	241	240	315	325	323	324	240
004	360	323	324	301	322	324	240	301	307
004	370	301	311	316	000	215	212	311	240
005	000	303	317	316	303	305	304	305	241

	005	010	000	215	212	311	240	315	317	326
	005	020	305	240	324	317	240	240	254	215
	005	030	212	311	240	327	311	316	241	240
	005	040	331	317	325	240	314	317	323	305
	005	050	241	000	215	212	304	322	301	327
	005	060	254	316	317	240	317	316	305	240
	005	070	327	311	316	323	256	000	215	212
	005	100	⇒310	305	322	305	247	323	240	324
	005	110	310	305	240	302	317	301	322	304
	005	120	215	212	261	336	262	336	263	215
	005	130	212	264	336	265	336	266	215	212
	005	140	267	336	270	336	271	215	212	311
	005	150	247	315	240	330	254	240	331	317
	005	160	325	247	322	305	240	317	215	212
	005	170	315	317	326	305	240	317	316	305
	005	200	240	323	321	325	301	322	305	240
	005	210	306	317	322	327	301	322	304	240
	005	220	311	306	240	326	301	303	301	316
	005	230	324	215	212	317	322	240	317	316
	005	240	305	240	323	321	325	301	322	305
ż	005	250	240	304	311	301	307	317	316	301
	005	260	314	314	331	240	324	317	240	303
	005	270	301	320	324	325	322	305	215	212
	005	300	331	317	325	240	323	324	301	322
	005	310	324	256	215	212	000	215	212	331
	005	320	317	325	240	327	311	316	256	000
	005	330	215	212	311	240	327	311	316	241
	005	340	240	331	317	325	240	310	301	326
	005	350	305	240	316	317	240	320	301	327
	005	360	316	323	256	000	215	212	316	317
	005	370	241	240	316	317	241	000		

006 000	000	INPUT
006 100	000	PRINT

After loading the Hexpawn program into memory, the program execution is begun by jumping to the start address of the program which is at location 000 on page 01. The program will print the introductory message followed by the starting position of the playing board. When the FM is displayed, the player enters the number of the square from which the pawn is to be moved. The program then prints TO, and the player enters the number of the square to which the pawn is to go. The program then makes its move, and the new board configuration is displayed. When the outcome of the game is evident to the program, a message is printed to indicate win, lose, or draw. A sample of three consecutive games is listed below. Note how the program goes through its learning process as the human player makes the same sequence of moves in each game.

> HERE'S THE BOARD 11213 41516 71819 I'M X, YOU'RE O MOVE ONE SQUARE FORWARD IF VACANT OR ONE SQUARE DIAGONALLY TO CAPTURE YOU START. XIXIX 1 1 01010 FM 8 TO 5 IXIX XIOL 01 10 FM 9 TO 6 XIOIX 01 FM 5 TO 2 YOU WIN.

XIXIX 1 1 01010 FM 8 TO 5 IXIX XIOI 0 10 FM 9 TO 6 IX X X O 011 FM 7 TO 5 11 XIOIX 11 FM 5 TO 2 YOU WIN. XIXIX 1 1 0000 FM 8 TO 5 IXIX XOI 0 10 FM 9 TO 6 X XIXIO FM 7 TO 5 I MOVE TO 7, I WIN! YOU LOSE!

AN 8080 LISTING OF THE HEXPAWN PROGRAM

This final listing is the operating portion of the Hexpawn program written for an 8080 based system. The 8080 version makes use of the more powerful instruction set of the 8080 mainly in setting up pointers, and includes instructions to set up the stack pointer, a function not required by the 8008. The operating portion of the 8080 version resides on pages 01 and 02 with the stack beginning at location 377 on page 02. The tables and messages are located on pages 03, 04, and 05 exactly as defined for the 8008 version. The user defined I/O routines should be set up as defined previously. The functional operation of the program is exactly as described in the text, and, therefore, need not be expanded upon. So, for those readers with 8080 based systems, here is the listing for the Hexpawn program.

	000		000		START,	LXS 000 003
	003		076	005		LXH 076 005
	006		165	002		CAL MSG
001	011		000			LXH 000 004
001	014	021	200	004		$LXD \ 200 \ 004$
001	017	032			RSTR,	LDAD
001	020	167				LMA
001	021	054				INL
001	022	034				INE
001	023	302	017	001		JFZ RSTR
001	026	041	000	003	AGAIN,	LXH 000 003
001	031	066	340			LMI 340
001	033	054				INL
001	034	066	007			LMI 007
001	036	056	302		PBD,	LLI 302
001	040	006	240			LBI 240
001	042	160				LMB
001	043	056	304			LLI 304
001	045	160				LMB
001	046	056	306			LLI 306
001	050	160				LMB
001	051	056	311			LLI 311
001	053	160				LMB
001	054	056	313			LLI 313
001	056	160				LMB
001	057	056	315			LLI 315
001	061	160				LMB
001	062	056	320			LLI 320
	064	160				LMB
001	065		322			LLI 322

001	067	160		LMB
001	070	$056 \ 324$		LLI 324
001	072	160		LMB
	073	056 000		LLI 000
	075	176		LAM
		054		INL
001	077	106		LBM
001	100	056 302		LLI 302
001	102	315 153	002	CAL STX
001	105	056 304		LLI 304
001	107	315 153	002	CAL STX
001	112	056 306		LLI 306
001	114	315 153	002	CAL STX
001	117	056 311		LLI 311
001	121	315 153	002	CAL STX
001	124	117		LCA
001	125	170		LAB
001	126	007		RLC
001	127	007		RLC
001	130	315 160	002	CAL STO
001	133	107		LBA
001	134	056 313	1	LLI 313
001	136	171		LAC
001	137	315 153	002	CAL STX
001	142	117		LCA
001	143	170		LAB
001	144	315 160	002	CAL STO
001	147	107		LBA
001	150	056 315		LLI 315
001	152	171		LAC
001	153	315 153	002	CAL STX
001	156	170		LAB
001	157	315 160	002	CAL STO
001	162	056 320	1	LLI 320
001	164	315 160	002	CAL STO
001	167	056 322	1	LLI 322
001	171	315 160	002	CAL STO
001	174	056 324	:	LLI 324
001	176	315 160	002	CAL STO
001	201	056 300		LLI 300

	001	203	315	165	002		CAL MSG
	001	206	056	002			LLI 002
	001	210	315	000	006		CAL INPUT
	001	213	167				LMA
	001	214	315	270	002		CAL FNUM
	001	217		177			JTS ERROR
	001		176				LAM
		223		017			NDI 017
	001	225	167	-			LMA
	001		107				LBA
	001		055				DCL
		230	176				LAM
	001			212	002		CAL RTAL
		234		177			JFC ERROR
		237		333			LLI 333
		241		165	002		CAL MSG
	001	244	056	003			LLI 003
	001	246	315	000	006		CAL INPUT
	001	251	167				LMA
	001	252	315	270	002		CAL FNUM
	001	255	372	177	002		JTS ERROR
	001	260	176				LAM
	001	261	346	017			NDI 017
	001	263	167				LMA
	001	264	055				DCL
	001	265	176				LAM
	001	266	326	003			SUI 003
	001	270	054				INL
	001		276				CPM
	001	272	312	231	002		JTZ BLK
	001	275	306	001			ADI 001
	001	277	276				CPM
	001	300	312	316	001		JTZ CKCAP
	001	303	326	002			SUI 002
ţ.	001	305	376	003			CPI 003
	001	307	312	177	002		JTZ ERROR
	001	312	276				CPM
	001	313	302	177	002		JFZ ERROR
	001	316		007		CKCAP,	
	001	320	312	177	002	,	JTZ ERROR

001	323	106				LBM
001	324	076	200			LAI 200
	326		223	002		CAL RTLP
	331	137				LEA
001	332	056	000			LLI 000
001	334	246				NDM
	335	312	177	002		JTZ ERROR
001	340	056			HMV,	LLI 002
	342	176				LAM
	343		220	002		CAL RTAR
	346	127				LDA
001	347	054				INL
001	350	176				LAM
001	351	376	004			CPI 004
			307	002		JTC HWIN
001	353 356	315	220	002		CAL RTAR
001		117				LCA
001	362	056	001			LLI 001
001	364	176				LAM
001	365	252				XRD
		$\begin{array}{c} 252 \\ 261 \end{array}$				ORC
001		117				LCA
001		167				LMA
001	371	055				DCL
001	372 373	173				LAE
001	373	247				NDA
100	374	312	001	002		JTZ NOCP
001	377	256				XRM
002		167				LMA
002	001	126			NOCP,	LDM
002	002	056	010			LLI 010
		172			SMDL,	LAD
002		276				CPM
002			032	002		JTZ OHLF
002		054				INL
002		054			SMD1,	INL
002	013	175				LAL
002	014	376				CPI 112
002	016	302	004			JFZ SMDL
002	021	041	340	004		LXH 340 004

002	024	315	165	002	CMSG,	CAL MSG
002	027	303	026	001		JMP AGAIN
002	032					
002	032	054			OHLF,	INL
002	033	171				LAC
002	034	276				CPM
002	035	302	012	002		JFZ SMD1
002	040					
002	040	055				DCL
002	041	175				LAL
002	042	017				RRC
002	043	306	106			ADI 106
002	045	157				LLA
002	046	156				LLM
002	047	046	004			LHI 004
	051	176			MFD1,	LAM
	052	247				NDA
			301			JTS ONO
002	056	302	065	002		JFZ MOVE
	061	054				INL
	062	303	051	002		JMP MFD1
	065					
	065	135			MOVE,	LEL
	066	041	004	003		LXH 004 003
	071	163				LME
002	072	007				RLC
	073	007				RLC
002	074	306	174			ADI 174
	076	157				LLA
	077	126				LDM
	100	054				INL
	101	116				LCM
002	102	054				INL
	103	136				LEM
	104	054				INL
	105	176				LAM
002	106	247				NDA
	107					JTS WIN
	112		262	002		JFZ DRAW
002	115	056	000			LLI 000

	117	176				LAM
	120	252				XRD
	121	261				ORC
	122	167				LMA
	123	054				INL
	124	173				LAE
	125	247				NDA
	126	312	036	001		JTZ PBD
002	131	256				XRM
002	132	167				LMA
002	133	302	036	001		JFZ PBD
002	136					
002	136	056	004			LLI 004
002	140	156				LLM
002	141	046	004			LHI 004
002	143	066	000			LMI 000
002	145	041	330	005		LXH 330 005
002	150	315	024	002		CAL CMSG
002	153					
002	153	007			STX,	RLC
002	154	320			,	RFC
002	155	066	330			LMI 330
002	157	311				RET
002	160					
002	160	007			STO,	RLC
002	161	320			,	RFC
002	162	066	317			LMI 317
002	164	311				RET
002	165					
002	165	176			MSG,	LAM
002	166	247			,	NDA
002	167	310				RTZ
002	170	315	100	006		CAL PRINT
	173	043				INXH
	174		165	002		JMP MSG
	177			· _		
	177	041	364	005	ERROR,	LXH 364 005
	202		165			CAL MSG
	205		003			LHI 003
	207		036	001		JMP PBD
004	_~.	500	500	UUT .		OTHE A DIV

	212	005				RTAL,	DCB
	213	310					RTZ
	214	007					RLC
	215	303	212	002			JMP RTAL
	220						,
	220	107				RTAR,	LBA
	221	076	001				LAI 001
	223	005				RTLP,	DCB
	224	310					RTZ .
	225	017					RRC
	226	303	223	002			JMP RTLP
	231						
002	231	106				BLK,	LBM
002	232	056	000				LLI 000
002	234	176					LAM
002	235	315	212	002			CAL RTAL
	240	247				SET,	NDA
002	241		177	002			JTS ERROR
002	244	036	000				LEI 000
002	246		340	001			JMP HMV
002	251						
002	251	062	025	005		WIN,	$\mathrm{STA}~025~005$
002	254	041	011	005			LXH 011 005
002	257	315	024	002			CAL CMSG
002	262						
002	262	041	052	005	· I	ORAW,	LXH 052 005
002	265	303	024	002			JMP CMSG
002	270						
	270	176]	FNUM,	LAM
002	271	376	260				CPI 260
002	273	370					RTS
002	274	326	272				SUI 272
002	276	306	200				ADI 200
002	300	311					RET
002	301						
002	301	041	374	004		ONO,	LXH 374 004
002	304	315	165	002		,	CAL MSG
002	307						
002	307	041	004	003		HWIN,	LXH 004 003
002	312	156					LLM

002 002 002 002 002 002 002	315 317 322 325	046 004 066 000 041 315 005 315 165 002 303 026 001		LHI 004 LMI 000 LXH 315 005 CAL MSG JMP AGAIN
006	000	000	INPUT,	
006	100	000	PRINT,	

HANGMAN!

HANGMAN is a word game with which most readers are probably well acquainted. The object of the game is to determine what word a player is thinking of by guessing the letters that make up the word. When characters contained in the word are correctly identified, the positions of the letters that have been ascertained are disclosed. The goal of the game is to ascertain all the letters making up the concealed word with the least amount of incorrect guesses. The game in the form to be presented traditionally received its name from the practice of creating a sketch of a stick figure being hung from a hangman's scaffold. A portion of the stick figure, such as a head, arms, torso, or legs, would be drawn in each time an incorrect letter guess was made. If the stick figure was completed before the entire word had been correctly identified, the player lost.

In the computerized version of the game to be presented here, the computer will select a word from a list of words (which may be created by the reader if desired). The computer will then allow a player to enter guesses as to the letters contained in the word selected. Each time the player correctly identifies a letter contained in the word, the characters that have been ascertained will be displayed in their proper location within the word. Each time a guess is incorrect, the computer will add a letter towards the spelling of Hangman! A game is finished when a player correctly identifies the word selected by the computer. Or, when eight incorrect letter guesses result in the complete spelling of Hangman!

The game is relatively simple to implement on a computer. However, despite its relative simplicity programming wise, the game can be surprisingly fun and challenging. This is due primarily to the nature of the game, augmented by the fact that the programmer has a virtually unlimited reservoir of alternatives to use when creating a list of words for the computer to select from when playing a game.

Besides its use as a pure fun game, the program can also be applied to more serious considerations such as making it a learning or teaching tool. Since the level of the vocabulary that is placed in the

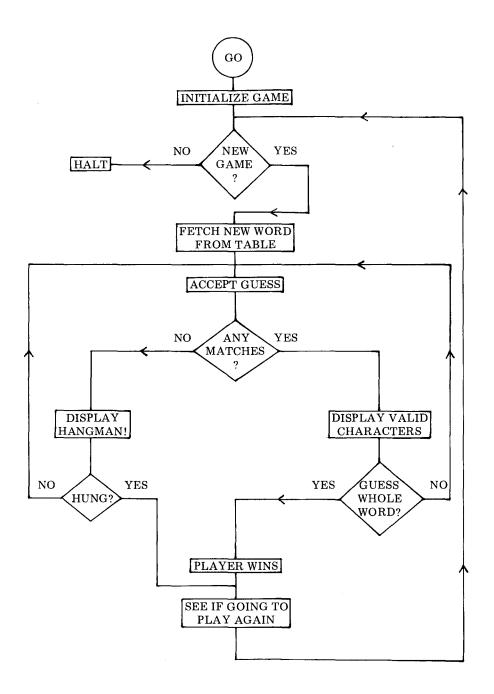
computer memory may be set as desired by the programmer, the game can be applied towards helping students develop vocabularies in virtually any subject. Additionally, one may readily change the language with which the game is played! French, Spanish, German, Maylasian! The computer will not care at all! The human player, though, may be suitably impressed with such variations!

FUNDAMENTAL STRUCTURE OF THE PROGRAM

The structure of the program is straightforward. Essentially, the computer is directed to select a word from a list of words in memory. A selected word is transferred to a buffer storage area. The player is asked to guess the letters in the selected word. Each time the player makes an entry, the buffer is scanned for any matches with the letter entered by the player. Appropriate matches are transferred to a working buffer that keeps track of all correct letter entries made by the player. Correct entries result in the contents of the working buffer being displayed to show the correct locations of letters properly identified by the player. Incorrect guesses by the player result in successive portions of the dreaded HANGMAN! being displayed. The overall flow of the program to be described here is illustrated in the flow chart on the following page.

DETAILS OF THE PROGRAM

The operating portion of the program described here fits easily into less than 1 K of memory excluding a variable length word table. The word table is simply a list of words. The list may extend as far as the user desires in available memory. A sample word table is included in this article. However, the reader may create a new list of words for the game. The word list provided uses about four pages of memory if the entire list is used. However, the list may be shortened if memory space is at a premium. A version of the program assembled to reside on pages 02 through 04 (operating portion) with the word table starting on page 05 will be provided as part of this article.



The first several routines in the program are used to initialize pointer storage locations and display a WANT A NEW WORD message to the system operator. Messages to be displayed by the program are stored as text strings in memory terminated by a zero byte. Text strings are displayed by calling a subroutine labeled MSG. MSG will output a string of characters pointed to by the H and L registers until a zero byte is detected. The actual MSG subroutine will be presented later. Suffice it to say at this point that one need only set up the H & L CPU registers to the starting address of a text string stored in memory, then call the MSG subroutine when it is desired to display such messages.

Following the WANT A NEW WORD message display, the program waits for a response from the operator by calling a user defined input subroutine labeled INPUTN. INPUTN should be designed by the reader to accept a character from the system's imput device (such as a keyboard) and return the character in the accumulator to the calling program. The INPUTN subroutine should also perform an echo display function so that the operator may verify the input character. The subroutine is free to utilize CPU registers A through E as far as this program is concerned. The user should note that this program expects the eighth bit in the accumulator to be in a '1' condition when the remaining seven bits represent an ASCII encoded character.

If the operator responds to the WANT A NEW WORD query by entering the letter N for no, the program terminates after displaying an appropriate response. If a Y for yes is entered, the program continues by calling upon a subroutine called MOVTAB. This subroutine will fetch a word from the program's word table. It will then transfer the word into a buffer. The buffer it is transferred into will be referred to as the word buffer in this article. The actual operation of the MOVTAB subroutine will be presented later. Suffice it to note at this time that upon return from the MOVTAB subroutine, a new word will be residing in the word buffer. The program will then be ready to start the playing of a game of Hangman!

START,	LHI 003	Set pointer to
	LLI 350	Number of guesses counter
	LMI 001	Initialize counter

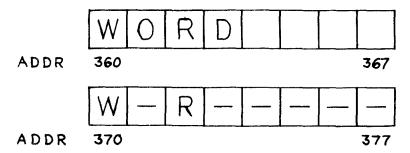
	LLI 356	Pointer to word table pointer
	LMI 000	Initialize to
	INL	Start of
	LMI 005	Word table
NEWONE,	LHI 004	Pointer to WANT A
· ·	LLI 000	NEW WORD? message
	CAL MSG	Display message
	CAL INPUTN	Fetch answer
	CPI 316	Was it NO?
	JTZ NOMORE	Say GOODBYE if no
	CPI 331	Else was it a YES?
	JFZ NEWONE	Ignore input otherwise
	CAL MOVTAB	If Y, fetch a word
	JMP GUESS	Go play the game
NOMORE,	LHI 004	Pointer to GOODBYE
,	LLI 025	Message
	CAL MSG	Display message
	HLT	End of playing session
		T - D - D

The next portion of the program begins by sending a message telling the operator to GUESS A LETTER. The program then accepts an input from the player. The input is expected to be any alphabetical character.

The program will then scan the word buffer to see if the letter received from the player matches with any of the letters in the word. Whenever a match is detected, the letter is stored in the same position in a second buffer called the guesses buffer.

It should be pointed out that the word buffer and the guesses buffer are identical in length. (Eight bytes in this program.) In the example program, the word buffer starts at location 360 on page 03. The guesses buffer starts at location 370 on page 03. At the start of each game, the word buffer, as previously mentioned, will be loaded with a word taken from the word table. A word may be up to eight letters in length. If there are not eight letters in a word, the balance of the word buffer will be filled with zero bytes. Furthermore, at the time the word buffer obtains a new word, the guesses buffer is filled with hyphens.

It thus becomes an easy matter to keep a record of correct guesses as the Hangman game progresses. Each time a position in the word buffer matches with the letter guessed by the player, the identical position in the guesses buffer is changed from a hyphen to the actual letter! The addressing scheme used in the program makes it easy to accomplish the objective. When a match is found in the word buffer, it is only necessary to add 010 (octal) to the buffer pointer to reach the corresponding position in the guesses buffer. The pictorial below should clarify the relationship.



If there is a match between the character inputted and any position in the word buffer, a flag is set (using register B). The word in the word buffer is scanned for a matching character until a zero byte is detected or the buffer pointer reaches the last address allocated for the buffer.

GUESS,	LHI 004 LLI 037 CAL MSG CAL INPUTN LCA	Pointer to GUESS A LETTER message Display message Fetch a letter Save letter in C
SCAN,	LBI 000 LHI 003 LLI 360	Clear B for a flag register Set pointer to Word buffer
СКМТСН,	CPM JFZ NOMTCH INB	Look for a match Skip ahead if not a match Set B as a flag

	LAI 010	Advance pointer
	ADL	To the guesses
	LLA	Buffer
	LMC	And deposit character
	LAL	Decrease pointer
	SUI 010	Back to
	LLA	Word buffer
NOMTCH,	INL	Advance buffer pointer
	LAM	See if next character
	NDA	Is a zero byte
	JTZ EOWORD	End of word if so
	LAI 007	See if at
	NDL	End of word buffer
	JTZ EOWORD	End of word if so
	LAC	Restore character to ACC
	JMP CKMTCH	Check next position

When the entire word buffer has been searched, the program checks the flag mentioned previously (in CPU register B) to determine if the player had made a correct letter guess. The flag will be set (have a value) if such was the case. It will still be zero if no match was detected.

If the flag was set during the SCAN operation, then the player has correctly determined a letter that exists in the word that the player is trying to identify. The program must now show the player how much of the word has been correctly identified. Thus, the program will first display an encouraging message. Then, the routine simply outputs the contents of the guesses buffer. The guesses buffer will contain all the locations of the letters that have been correctly identified during the game. Unidentified locations will still contain a hyphen. Thus, if a player had correctly identified the letters W and R in the spelling of WORD, the computer would output: W-R-.

As the program outputs the contents of the guesses buffer, it checks to see if any hyphens (referred to in the listing as dashes) are displayed. A software flag mechanism is used for this purpose. At the end of the guesses buffer outputting operation, the flag is tested. If it is zero, then the player has identified the entire word. The program will then display a congratulations message and go back to see if the player wants to continue the game with a new word. If there are any hyphens left in the guesses buffer as indicated by the software flag being set, then the program loops back to allow the player to guess another letter.

The reader may note that the routine examines the word buffer to determine when to stop outputting the contents of the guesses buffer. This is because the word buffer will contain a zero byte at the end of the word if the word is less than eight characters in length. The guesses buffer does not contain such an indicator.

The portion of the program just discussed is presented next.

EOWORD,	INB DCB JTZ HANGIT LHI 004 LLI 074 CAL MSG LHI 003 LLI 353 LMI 000 LLI 370	At end of word, exercise The MATCH flag If = 0, no matches If match(es), set pointer To GOOD. YOU HAVE: msg Display message Set pointer to Dashes counter storage Clear dashes counter Pointer to guesses buffer
NOTEND,	LAM CPI 255 JFZ AHEAD2 LEL LLI 353 LBM INB LMB LLE	Fetch a character See if it is a dash Skip next instruction if not Save pointer temporarily Set pointer to dashes counter Fetch dashes value Increment Restore to memory Restore saved pointer
AHEAD2,	CAL PRINT INL LAL SUI 010	Print the character Advance the buffer pointer Decrease pointer To word buffer

	LLA LAM NDA JTZ ENDAGN LAL NDI 007 JTZ ENDAGN LAI 010 ADL LLA JMP NOTEND	Here Fetch data from word buffer And see if it is Zero byte, jump if so Fetch pointer See if at end of word Buffer, jump if so Else restore pointer Ahead to The guesses buffer Do next character in buffer
ENDAGN,	LLI 353 LBM INB DCB JFZ GUESS LHI 004 LLI 120 CAL MSG JMP NEWONE	Pointer to dashes counter Fetch value Exercise the dashes Flag register Word not completed If reach here, set pointer To congratulations message Display message Go play with a new word

For the case when the player has inputted a letter that does not exist in the word in the word buffer, the program must take a different course of action. This case is handled by a portion of the program that starts at the label HANGIT. Here the operator is informed of the incorrect guess by the display of the message NOPE. This message is then followed by the display of a portion or all of the HANGMAN! message.

Each time the player guesses incorrectly during a game, a letter is added to the message spelling out the word HANGMAN! In order to do this properly, the program maintains a counter of the number of incorrect guesses made. Then, the computer is simply used to determine how many letters of the HANGMAN! message to display. (The HANGMAN! message is simply stored in a buffer, starting at location 340 on page 03 in the example program.) If the entire HANGMAN! statement is not displayed, the balance of the message is shown as dash signs (hyphens). The number of dash signs to display is calculated by subtracting the value of the counter (of incorrect guesses made) from 10 (octal) which is the number of characters in the HANGMAN! message. Thus, as the game progresses, the message HANGMAN! will appear more and more complete with each incorrect guess as illustrated below.

H HA	(First incorrect guess) (Second incorrect guess)
HAN	(Third incorrect guess)
HANGMAN!	(Eighth incorrect guess)

When eight incorrect guesses have been made in a game, the entire HANGMAN! message will be displayed. The player then loses the game. The program will then go back and see if the player wants to start with a new word.

The listing for the portion of the program that displays the dreaded HANGMAN! message is shown next.

HANGIT,	LHI 004 LLI 062 CAL MSG LHI 003 LLI 350 LBM INB LMB INL LMB LAI 010 SUB INL LMA LEI 340	Pointer to NOPE Message Display message Set pointer to Number of guesses in counter Fetch counter Increment it Restore the counter Advance pointer Save it again Calculate number of Dashes left in HANGMAN! Advance pointer Save the value in memory Init. pntr. to HANGMAN bfr
HANGMR,	LLE LAM CAL PRINT	Set pntr to HANGMAN bfr Fetch a character from buffer Display it

	INL LEL LLI 351 LBM DCB LMB JFZ HANGMR LLI 352 LCM INC DCC JTZ NEWONE LAI 255	Advance the buffer pointer Save temporarily in E Pointer to counter in memory Fetch counter value Decrement Restore to memory Continue if counter not zero Pntr to second cntr in memory Fetch counter value Exercise counter value To see if it is zero Start new game if so Else load code for "-"
MRDASH,	CAL PRINT LCM DCC LMC JFZ MRDASH JMP GUESS	Display a dash Fetch counter Decrement Restore Until counter is zero Then continue game

The next portion of the program is a subroutine mentioned previously called MOVTAB. The primary function of this subroutine is to fetch a new word from a list or table of words stored in memory. However, the subroutine also performs a few other functions that need to be performed each time a new word (actually representing the start of a new game) is obtained.

The first thing the subroutine does is fetch the value of the counter used for keeping track of how many incorrect guesses were made during the last game played. This value will be used to determine how many words in the word table to skip over when selecting a new word. This method is used so that words will be selected from the table in a rather arbitrary fashion rather than simply taking the next word in the list. If the next word in the list was always taken, players might soon start remembering certain words or sequences of words which would soon make the game somewhat boring!

The program then initializes the guesses buffer to the all hyphens

condition by loading the ASCII code for the dash sign (255 octal) into all the locations in the buffer. In a similar fashion, the word buffer is cleared to the all zeroes condition in preparation for its receiving a new word from the word table.

These initial functions of the subroutine are shown below.

MOVTAB,	LHI 003 LLI 350 LBM LLI 370 LCI 010 LAI 255	Pointer to number Of guesses counter Fetch and save in B Pointer to guesses buffer Set a loop counter Set code for "-"
DASHFL,	LMA INL DCC JFZ DASHFL	Fill guesses buffer With dashes Until counter Is zero
NXWORD,	LLI 360 LCI 010 XRA	Set pointer to word buffer Set a loop counter Clear the accumulator
ZEROFL,	LMA INL DCC JFZ ZEROFL	Fill word buffer With zero bytes Until counter Is zero

Before explaining the operation of the portion of the subroutine that extracts a new word from the word table, it will be beneficial to explain the organization of the table.

The table consists of a list of words stored in memory in the following format.

А	1st letter of a word
A+1	2nd letter of a word
A+2	3rd letter of a word

•	•
A+N	Nth letter of a word
A+N+1	000 word terminator code
В	1st letter of a word
B+1	2nd letter of a word
•	
B+N	Nth letter of a word
B+N+1	000 word terminator code
\mathbf{C}	1st letter of a word
•	
	•
C+N+1	000 word terminator code
D	000 end of table terminator

The reader should notice that each word in the list must be terminated by a zero byte. Words must also be limited to eight or less letters in length (or they would overflow the word buffer). The table is terminated by placing an additional zero byte immediately following the zero byte word terminator after the last word in the table. The word table in the program provided in this manual starts on page 05 at location 000. The table may extend for as long as the user desires within available memory. (In the sample word list about 100 words are provided. These require about three pages of memory. Of course, the list may be shortened if necessary. Or the user may provide a completely original table of words.)

A word is extracted from the word table through the following procedure. First, a word table pointer is extracted from its storage location in memory and loaded into CPU registers H and L. This pointer will initially point to the first letter of a word in the table. Next, a character is extracted from the word table. The character obtained is first tested to see if it is a zero byte. A zero byte in place of an expected letter (as the first letter in a word) indicates that the end of the table has been reached. In that case, the pointer in H and L is reset back to the start of the word table.

Next, a second pointer is established in CPU registers D and E.

This pointer will be used to point to the word buffer during transfer operations from the word table. Now a character is fetched from the word table. Then the pointers in H and L and D and E are swapped and the character is transferred into the word buffer. (Unless a zero byte indicating the end of a word is detected. In that case no transfer takes place.) Next, the two sets of pointers are advanced. The process is then repeated until a whole word has been loaded into the word buffer.

When an entire word has been transferred into the word buffer, the routine advances the word table pointer once more. (This is so it will be advanced over the end of word terminator and be pointing at the first letter in the next word in the table.) Then the pointer is restored to its storage location in memory. Next, the routine fetches the number of guesses counter. It decrements the value of that counter. If the value is not zero after the decrement operation, then the routine loops back (to the label NXWORD), and proceeds to read the next word in the word table into the word buffer. (The reason for following this procedure was presented earlier.) When the counter reaches zero, it is stored in memory (at its 000 value) for use during the next game. The subroutine is then exited.

	LLI 356 LAM INL LHM LLA XRA CPM JFZ AHEAD3 LHI 005 LLI 000	Set pointer to word table pntr Fetch the low address Advance pointer Set the page address And low address Clear the accumulator See if first entry is zero Skip ahead if not Reset pointer to start Of word table if so
AHEAD3,	LDI 003 LEI 360	Set pointer to word Buffer in D and E
BUFFMR,	LAM NDA JTZ NEXT	Fetch a character from table Exercise flags If zero, have whole word

	CAL SWITCH LMA INL CAL SWITCH INL JFZ NOHIGH INH	Else swap pointers Dep character in word buffer Advance word buffer pointer Swap pointers Advance word table low pntr If not zero, skip next Advance table high pointer
NOHIGH,	JMP BUFFMR	Continue transfer from table
NEXT,	INL JFZ NOTHI INH	Advance table pointer Low address And high address if required
NOTHI,	CAL SWITCH LLI 356 LME INL LMD DCB JFZ NXWORD LLI 350 LMI 000 RET	Save pointer in D and E Set pointer to table pointer Save table pointer low Address and High address Decr number guesses counter If not zero, get next word Else set pointer to guesses Counter and zero counter Then exit subroutine

That completes the discussion of the major routines in the program. There are two more minor utility subroutines used in the program. One of these is simply a subroutine called SWITCH that is used to exchange the contents of CPU registers H and L with D and E. During this operation, CPU register C is used as a temporary register.

SWITCH,	LCH	Put H into C temporarily
	LHD	Load D into H
	LDC	Now orig H from C to D
	LCL	Put L into C temporarily
	LLE	Load E into L
	\mathbf{LEC}	Now orig L from C to E
	RET	Swapping oper. completed

The other is the subroutine mentioned earlier called MSG. MSG simply outputs a string of characters from memory to an output device until it detects a zero byte.

MSG,	LAM	Fetch a character
	NDA	See if a zero byte
	RTZ	Indicating end of string
	CAL PRINT	If not, display character
	INL	Increment low address pointer
	JFZ MSG	Get next character unless
	INH	Need to advance page address
	JMP MSG	Then get next character

The MSG subroutine above calls on another subroutine which has been termed PRINT. The PRINT subroutine must be an actual device operating subroutine that will cause the ASCII character in the accumulator to be transmitted to the output device being used by the system. The PRINT subroutine, which must be provided by the user, may use the CPU registers B through E if required. It should not alter the contents of the H and L CPU registers (unless the subroutine is able to restore those registers to their original values at the conclusion of the process).

ASSEMBLED LISTING OF THE "HANGMAN!" PROGRAM FOR AN '8008'

An assembled listing of the program for operation on an 8008 system is presented next. The operating portion of the program has been assembled to reside in pages 02 and 03. Page 04 is reserved for the various message strings used by the program plus the user provided I/O subroutines. The word table for the program is assumed to start on page 05. A sample list of words for use with the program is provided in ASCII form at the end of the assembled listing.

002 002 002 002 002	000 002 004 006 010 012 013	056003066350076001066356076000060005		START,	LHI 003 LLI 350 LMI 001 LLI 356 LMI 000 INL LMI 005
002 002 002 002 002 002 002 002	$\begin{array}{c} 015\\ 017\\ 021\\ 024\\ 027\\ 031\\ 034\\ 036\\ 041\\ 044 \end{array}$	$\begin{array}{ccccc} 056 & 004 \\ 066 & 000 \\ 106 & 110 \\ 106 & 200 \\ 074 & 316 \\ 150 & 047 \\ 074 & 331 \\ 110 & 015 \\ 106 & 347 \\ 104 & 057 \end{array}$	003 004 002 002 002	NEWONE,	LHI 004 LLI 000 CAL MSG CAL INPUTN CPI 316 JTZ NOMORE CPI 331 JFZ NEWONE CAL MOVTAB JMP GUESS
002 002 002	047 051 053 256	$\begin{array}{c} 104 & 001 \\ 056 & 004 \\ 066 & 025 \\ 106 & 110 \\ 000 \end{array}$		NOMORE,	LHI 004 LLI 025 CAL MSG HLT
002 002 002	057 061 063 066 071	056 004 066 037 106 110 106 200 320	003	GUESS,	LHI 004 LLI 037 CAL MSG CAL INPUTN LCA
002	$\begin{array}{c} 072 \\ 074 \\ 076 \end{array}$	016 000 056 003 066 360		SCAN,	LBI 000 LHI 003 LLI 360
002 002 002 002 002	100 101 104 105 107 110 111	$\begin{array}{c} 277\\ 110\\ 010\\ 006\\ 010\\ 206\\ 360\\ 372 \end{array}$		СКМТСН,	CPM JFZ NOMTCH INB LAI 010 ADL LLA LMC

002	112	306				LAL
002	113	024	010			SUI 010
002	115	360				LLA
002	116	060			NOMTCH,	INL
002	117	307				LAM
002	120	240				NDA
002	121	150	136	002		JTZ EOWORD
002	124	006	007			LAI 007
002	126	246				NDL
002	127	150	136	002		JTZ EOWORD
002	132	302				LAC
002	133	104	100	002		JMP CKMTCH
002	136	010			EOWORD,	INB
002	137	011				DCB
002	140	150	253	002		JTZ HANGIT
002	143	056	004			LHI 004
002	145	066	074			LLI 074
002	147	106	110	003		CAL MSG
002	152	056	003			LHI 003
002	154	066	353			LLI 353
002	156	076	000			LMI 000
002	160	066	370			LLI 370
002	162	307			NOTEND,	LAM
002	163	074	255			CPI 255
002	165	110	177	002		JFZ AHEAD2
002	170	346				LEL
002	171	066	353			LLI 353
002	173	317				LBM
002	174	010				INB
002	175	371				LMB
002	176	364				LLE
	177	106	300	004	AHEAD2,	CAL PRINT
002		060			·	INL
002	203	306				LAL
	204	024	010			SUI 010
002	206	360				LLA

		207					LAM
	002	210	240				NDA
	002	211	150	231	002		JTZ ENDAGN
	002	214	306				LAL
		215		007	,		NDI 007
	002	217	150	231	002		JTZ ENDAGN
	002	$\begin{array}{c} 217 \\ 222 \end{array}$	006	010			LAI 010
	002	224	206				ADL
	002	225	360				LLA
	002	226	104	162	002		JMP NOTEND
	002	231	066	353		ENDAGN,	LLI 353
	002	$\begin{array}{c} 231 \\ 233 \end{array}$	317				LBM
	002	234	010				INB
	002	235	011				DCB
	002	236	110	057	002		JFZ GUESS
	002	$\begin{array}{c} 241 \\ 243 \end{array}$	056	004			LHI 004
	002	243	066	120			LLI 120
	002	245	106	110	003		CAL MSG
	002	250	104	015	002		JMP NEWONE
	002	253	056	004		HANGIT,	LHI 004
	002		066				LLI 062
	002	257	106	110	003		CAL MSG
		262					LHI 003
	002	264	066	350			LLI 350
			317				LBM
•	002	267	010				INB
	002	$\begin{array}{c} 270\\ 271 \end{array}$	371				LMB
	002	271	060				INL
	002		371				LMB
	002	273	006	010			LAI 010
	002	275	221				SUB
	002	276	060				INL
	002	277	370				LMA
	002	300	046	340			LEI 340
	002	302	364			HANGMR,	LLE
		303					LAM
	002	304	106	300	004		CAL PRINT

002 002 002 002 002 002 002 002 002 002	307 310 311 313 314 315 316 321 323 324 325 326 331	$\begin{array}{c} 060\\ 346\\ 066\\ 3\\ 317\\ 011\\ 371\\ 110\\ 3\\ 066\\ 3\\ 327\\ 020\\ 021\\ 150\\ 006\\ 2\end{array}$	302 352 015			INL LEL LLI 351 LBM DCB LMB JFZ HANGMR LLI 352 LCM INC DCC JTZ NEWONE LAI 255
002 002 002 002	333 336 337 340 341 344	106 3 327 021 372 110 3 104 0	333	002	MRDASH,	CAL PRINT LCM DCC LMC JFZ MRDASH JMP GUESS
002 002 002 002	347 351 353 354 356 360	$\begin{array}{ccc} 056 & 0 \\ 066 & 3 \\ 317 \\ 066 & 3 \\ 026 & 0 \\ 006 & 2 \end{array}$	350 370 010		MOVTAB,	LHI 003 LLI 350 LBM LLI 370 LCI 010 LAI 255
$\begin{array}{c} 002\\ 002 \end{array}$	$362 \\ 363 \\ 364 \\ 365$	370 060 021 110 3	362	002	DASHFL,	LMA INL DCC JFZ DASHFL
002	$370 \\ 372 \\ 374$	$\begin{array}{ccc} 066 & 3 \\ 026 & 0 \\ 250 \end{array}$			NXWORD,	LLI 360 LCI 010 XRA
002	375 376 377	370 060 021			ZEROFL,	LMA INL DCC

	003 005 006 007 010 011 012 013 016	066 307 060 357 360 250 277	022 005			JFZ ZEROFL LLI 356 LAM INL LHM LLA XRA CPM JFZ AHEAD3 LHI 005 LLI 000
003 003	$\begin{array}{c} 022\\ 024 \end{array}$	036 046			AHEAD3,	LDI 003 LEI 360
003 003 003 003 003 003 003	$\begin{array}{c} 026\\ 027\\ 030\\ 033\\ 036\\ 037\\ 040\\ 043\\ 044\\ 047 \end{array}$	106 370 060 106	053 101 101 050	003 003	BUFFMR,	LAM NDA JTZ NEXT CAL SWITCH LMA INL CAL SWITCH INL JFZ NOHIGH INH
003	050	104	026	003	NOHIGH,	JMP BUFFMR
	053 054 057	$060 \\ 110 \\ 050$	060	003	NEXT,	INL JFZ NOTHI INH
	063 065 066 067 070 071	066 374 060 373 011	370		NOTHI,	CAL SWITCH LLI 356 LME INL LMD DCB JFZ NXWORD LLI 350

003 003		076 007	000			LMI 000 RET
003 003 003 003 003 003 003	102 103 104 105 106	325 353 332 326 364 342 007			SWITCH,	LCH LHD LDC LCL LLE LEC RET
003 003 003	$ \begin{array}{r} 111 \\ 112 \\ 113 \\ 116 \\ 117 \\ \end{array} $	060 110 050	300 110 110	003	MSG,	LAM NDA RTZ CAL PRINT INL JFZ MSG INH JMP MSG
003 003 003 003		310 301 316 307 315 301 316 241				310 301 316 307 315 301 316 241
$ \begin{array}{c} 003 \\ 003 \\ 003 \end{array} $	352 353	000 000 000 000				000 000 000 000
003	360 361	000 000 000 000 000				000 000 000 000 000

003	363	000							000	
003	364	000							000	
003	365	000							000	
003	366	000							000	
003	367	000							000	
003	370	255							255	
003	371	255							255	
003	372	255							255	
003	373	255							255	
003	374	255							255	
003	375	255							255	
003	376	255							255	
003	377	255							255	
	•••								200	
	004	4 000	215	212	212	327	301	316	324	240
	004		301	240	316	305	327	240	327	317
		4 020	322		277		000	215	212	307
	004		317	317	304		331	$\frac{1}{241}$	000	215
		4 040	212	307		305	323	323	2 40	301
	004		240	314	305	324	324	305	322	272
		4 060	240	000	215	212	316	317	320	305
	004		$\frac{10}{241}$	240	240	000	215	212	307	317
	004		317	304	256	240	331	317	325	240
	004		310	301	326	305	272	240	240	000
	004		215	212	303	317		307	322	301
	004		304	325	314	301	324	311	317	316
	00-		323	$\frac{525}{241}$	000	001	024	011	017	910
	00-	1 140	040	441 1	000					

004 200

INPUTN,

004 300

PRINT,

005	000	310	305	314	314	317	000	301	322
005	010	322	317	327	000	303	317	315	320
005	020	325	324	305	322	000	320	322	305
005	030	315	311	325	315	000	316	317	324
005	040	311	303	305	000	306	325	316	000

005	050	310	305	301	326	331	000	322	325
005	060	302	302	305	322	000	322	325	323
005	070	324	314	305	000	324	310	327	301
005	100	322	324	000	317	331	323	324	305
005	110	322	000	317	330	311	304	311	332
005	120	305	000	317	323	323	311	306	331
005	130	000	317	320	311	316	311	317	316
005	140	000	317	317	332	331	000	317	316
005	150	305	322	317	325	323	000	316	317
005	160	315	301	304	000	316	317	303	324
005	170	325	322	316	305	000	316	317	315
005	200	311	316	301	324	305	000	316	325
005	210	315	323	313	325	314	314	000	304
005	220	301	306	306	317	304	311	314	000
005	230	323	311	304	305	322	305	301	314
005	240	000	303	322	311	303	313	305	324
005	250	000	303	317	325	322	311	305	322
005	260	000	303	317	323	315	317	323	000
005	270	303	310	305	315	311	323	324	000
005	300	303	310	305	315	311	303	301	314
005	310	000	303	310	311	303	317	322	331
005	320	000	303	310	314	317	322	311	316
005	330	305	000	303	311	324	311	332	305
005	340	316	000	303	311	324	322	325	323
005	350	000	303	314	317	323	305	324	000
005	360	303	317	307	305	316	324	000	302
005	270	311	322	304	000	302	305	305	324
006	000	314	305	000	302	305	314	311	305
006	010	326	305	000	302	301	324	310	324
006	020	325	302	000	302	301	323	313	305
006	030	324	000	302	301	316	321	325	305
006	040	324	000	302	301	302	302	311	324
006	050	324	000	302	301	303	313	302	317
006	060	316	305	000	301	325	304	311	302
006	070	314	305	000	301	323	320	311	322
006	100	311	316	000	301	323	324	305	322
006	110	317	311	304	000	301	320	320	322
006	120	317	326	301	314	000	301	320	317
006	130	307	305	305	000	301	316	316	325
006	140	311	324	331	000	301	316	317	304

	-								
006	150	311	332	305	000	301	314	325	315
006	160	311	316	325	315	000		311	322
006	170	000	301	311	323	314	305	000	301
006	200	304	312	317	311	316	000	301	302
006	210	331	323	323	000	301	302	317	314
006	220	311	323	310	000	321	325	305	325
006	230	305	000	321	325	311	326	305	322
006	240	000	321	325	301	314	315	000	321
006	250	325	311	324	305	000	321	325	311
006	260	330	317	324	311	303	000	321	325
006	270	317	311	316	000	321	325	317	311
006	300	324	000	321	325	317	324	311	305
006	310	316	324	000	322	301	304	311	317
006	320	000	322	301	311	323	311	316	000
006	330	322	301	320	324	000	322	301	324
006	340	311	317	000	322	301	325	303	317
006	350	325	323	000	322	301	331	317	316
006	360	000	322	301	332	317	322	000	322
006	370	305	301	314	315	000	322	305	305
007	000	313	000	322	305	307	311	323	324
007	010	305	322	000	322	311	326	305	324
007	020	000	323	303	310	317	317	316	305
007	030	322	000	323	301	325	316	301	000
007	040	323	301	324	311	316	000	323	303
007	050	305	320	324	305	322	000	323	303
007	060	311	305	316	303	305	000	323	303
007	070	322	311	302	302	314	305	000	302
007	100	305	310	311	316	304	000	304	311
007	110	307	316	311	306	331	000	305	314
007	120	314	311	320	324	311	303	000	305
007	130	314	317	321	325	305	316	324	000
007	140	305	314	325	323	311	326	305	000
007	150	306	305	301	324	310	305	322	000
007	160	307	301	314	314	317	327	323	000
007	170	307	301	322	304	305	316	000	307
007	200	301	332	305	314	314	305	000	315
007	210	301	303	301	302	322	305	000	326
007	220	301	314	311	301	316	324	000	326
007	230	305	316	311	323	317	316	000	326
007	240	311	326	311	304	000	327	305	311

007	250 260 270	304	000	327	000 311 000	323	305	000	332
007 007 007 007	300 310 320 330 340	317 324 331 314	307 310 317 317 317	331 000 314 327	000 331 313 000 311	332 301 000 331	305 327 331 325	316 316 305 314	311 000 314 305
~ ~ .	350				304			011	000

A list of the messages used in the game (which reside on page 04 in the assembled listing just presented) is shown below in the order in which they appear in the messages table.

WANT A NEW WORD? GOODBYE! GUESS A LETTER: NOPE! GOOD. YOU HAVE: CONGRATULATIONS!

For those that want to use the word list supplied as an example, (pages 05 through 07 in the listing just presented) the list on the following page will serve as a reference. The words appear in the same order as they are stored in the list. (Remember, however, that the program will skip around the list as it selects the next word that will be played!)

HELLO	BASKET	SAUNA
ARROW	BANQUET	SATIN
COMPUTER	BABBITT	SCEPTER
PREMIUM	BACKBONE	SCIENCE
NOTICE	AUDIBLE	SCRIBBLE
FUN	ASPIRIN	BEHIND
HEAVY	ASTEROID	DIGNIFY
RUBBER	APPROVAL	ELLIPTIC
RUSTLE	APOGEE	ELOQUENT
THWART	ANNUITY	ELUSIVE
OYSTER	ANODIZE	FEATHER
OXIDIZE	ALUMINUM	GALLOWS
OSSIFY	AIR	GARDEN
OPINION	AISLE	GAZELLE
OOZY	ADJOIN	MACABRE
ONEROUS	ABYSS	VALIANT
NOMAD	ABOLISH	VENISON
NOCTURNE	QUEUE	VIVID
NOMINATE	QUIVER	WEIGHT
NUMSKULL	QUALM	WEIRD
DAFFODIL	QUITE	WISE
SIDEREAL	QUIXOTIC	ZERO
CRICKET	QUOIN	ZOOLOGY
COURIER	QUOIT	ZENITH
COSMOS	QUOTIENT	YAWN
CHEMIST	RADIO	YOLK
CHEMICAL	RAISIN	YELLOW
CHICORY	RAPT	YULE
CHLORINE	RATIO	TRICKLE
CITIZEN	RAUCOUS	END
CITRUS	RAYON	
CLOSET	RAZOR	
COGENT	REALM	
BIRD	REEK	
BEETLE	REGISTER	
BELIEVE	RIVET	
BATHTUB	SCHOONER	

Once the program has been loaded into memory (along with the user provided I/O routines!) the program is ready to operate. Simply start program execution at page 02 location 000. Operation from then on is directed by the program. A sample of the program's operation is illustrated below.

WANT A NEW WORD? Y GUESS A LETTER: Α NOPE! H-----**GUESS A LETTER:** Ε GOOD. YOU HAVE: - E - - -GUESS A LETTER: R NOPE! HA-----GUESS A LETTER: L GOOD. YOU HAVE: -ELL-GUESS A LETTER: В NOPE! HAN----GUESS A LETTER: \mathbf{S} NOPE! HANG---**GUESS A LETTER:** 0 GOOD. YOU HAVE: -ELLO GUESS A LETTER: Н GOOD. YOU HAVE: HELLO **CONGRATULATIONS!** WANT A NEW WORD? Y GUESS A LETTER: Α NOPE! H-----GUESS A LETTER: E ----E GOOD. YOU HAVE: T

GUESS A LETTER: I GOOD. YOU HAVE: --- I - E GUESS A LETTER: O GOOD. YOU HAVE: - O - I - E GUESS A LETTER: U NOPE! HA-----

GUESS A LETTER:	R
NOPE! HAN	
GUESS A LETTER:	Т
GOOD. YOU HAVE:	- OTI - E
GUESS A LETTER:	Ν
GOOD. YOU HAVE:	NOTI - E
GUESS A LETTER:	С
GOOD. YOU HAVE:	NOTICE
CONGRATULATION	S!

WANT A NEW WORD? Y **GUESS A LETTER:** W NOPE! H-----GUESS A LETTER: Y NOPE! HA-----GUESS A LETTER: Р NOPE! HAN----GUESS A LETTER: Ν NOPE! HANG---**GUESS A LETTER:** G NOPE! HANGM --**GUESS A LETTER:** V NOPE! HANGMA-GUESS A LETTER: С NOPE! HANGMAN-**GUESS A LETTER:** Х NOPE! HANGMAN!

WANT A NEW WORD? N GOODBYE!

The program will continue to operate until a player responds with a N for NO to the WANT A NEW WORD query. At that time, the program will halt. If it is desired to continue playing after a NO response to that question, the program may simply be restarted at the starting address (page 02 location 000).

ASSEMBLED LISTING OF THE PROGRAM FOR AN 8080 SYSTEM

The following is an assembled listing of the HANGMAN! program designed to run on an 8080 system. Only minor changes have been made in the program to take advantage of some of the special capabilities of the 8080 instruction set. However, the basic organization of the program has not been altered so that the previous detailed discussion of the program's operation still applies. The message table and word list would be in the same format as the 8008 version. Of course, the user will need to provide the appropriate I/O routines for either version of the program. They can be placed in the same memory locations for the following 8080 version as was suggested for the 8008 example (on page 04 starting at locations 200 (input) and 300 (output)).

002 002 002	003 005 007	041 066 056 066 054 066	356 000	003	START,	LXH 350 003 LMI 001 LLI 356 LMI 000 INL LMI 005
002 002 002 002 002 002 002 002 002	014 017 022 025 030 032 035 037 042 045	041 315 315 376 312 376	316 050 331 014 327	004 003 004 002	NEWONE,	LXS 200 004 LXH 000 004 CAL MSG CAL INPUTN CPI 316 JTZ NOMORE CPI 331 JFZ NEWONE CAL MOVTAB JMP GUESS
002	050 053 056	$041 \\ 315 \\ 166$		004 003	NOMORE,	LXH 025 004 CAL MSG HLT

002		315	032	003	GUESS,	LXH 037 004 CAL MSG CAL INPUTN LCA
$\begin{array}{c} 002\\ 002 \end{array}$	071 073	006 041	000 360	003	SCAN,	LBI 000 LXH 360 003
002 002 002 002 002 002 002 002	102 103 105 106 107 110	004	114 010 010	002	СКМТСН,	CPM JFZ NOMTCH INB LAI 010 ADL LLA LMC LAL SUI 010 LLA
002 002 002 002	115 116 117 122 124 125 130	245 312 171		002	NOMTCH,	INL LAM NDA JTZ EOWORD LAI 007 NDL JTZ EOWORD LAC JMP CKMTCH
002 002 002 002 002 002	136	041 315 041 066	243 074 032 353 000 370	004 003	EOWORD,	INB DCB JTZ HANGIT LXH 074 004 CAL MSG LXH 353 003 LMI 000 LLI 370
	$\begin{array}{c} 156 \\ 157 \end{array}$	$\begin{array}{c} 176\\ 376 \end{array}$	255		NOTEND,	LAM CPI 255

002 002 002	161 164 165 167 170	302 135 056 064 153	171 353	002		JFZ AHEAD2 LEL LLI 353 INM LLE
002 002 002 002 002	171 174 175 176 200 201 202	315 054 175 326 157 176 247	300 010	004	AHEAD2,	CAL PRINT INL LAL SUI 010 LLA LAM NDA
002 002		247 312 175 346		002		NDA JTZ ENDAGN LAL NDI 007
002 002	$211 \\ 214 \\ 216 \\ 217$	312 076 205 157	223 010	002		JTZ ENDAGN LAI 010 ADL LLA
002 002 002 002 002	220 223 225 226 227	303 056 064 065 302	057	002	ENDAGN,	JMP NOTEND LLI 353 INM DCM JFZ GUESS
002	$232 \\ 235 \\ 240$	315	$120 \\ 032 \\ 014$	003		LXH 120 004 CAL MSG JMP NEWONE
002 002 002 002 002 002 002 002 002	$\begin{array}{c} 243\\ 246\\ 251\\ 254\\ 255\\ 256\\ 257\\ 260\\ 262\\ 263\\ 264\\ \end{array}$	315	062 032 350 010	003	HANGIT,	LXH 062 004 CAL MSG LXH 350 003 INM LAM INL LMA LAI 010 SUM INL LMA

002	265	036	340			LEI 340
	267 270 271	$153 \\ 176 \\ 315$	300	004	HANGMR,	LLE LAM CAL PRINT
002	274	054	000	001		INL
	$\begin{array}{c} 275\\ 276 \end{array}$	135 056	251			LEL LLI 351
	300	065	001			DCM
002			267	002		JFZ HANGMR
002	304 306	056 064	352			LLI 352 INM
002	307	065				DCM
	310			002		JTZ NEWONE
002	313	076	255			LAI 255
002	315		300	004	MRDASH,	CAL PRINT
002	320	065				DCM
	321 324		$\begin{array}{c} 315\\ 057 \end{array}$			JFZ MRDASH JMP GUESS
002	024	000	001	002		JMF GUESS
	327		350	003	MOVTAB,	LXH 350 003
	332 333	$\begin{array}{c} 106 \\ 056 \end{array}$	370			LBM LLI 370
	335	016				LCI 010
	337	076	255			LAI 255
002	341	167			DASHFL,	тма
002	342	054			DASHI L,	INL
	343	015				DCC
002	344	302	341	002		JFZ DASHFL
	347	041	360	003	NXWORD,	LXH 360 003
	352	016	010		,	LCI 010
002	354	257				XRA
002		167			ZEROFL,	LMA
		054				INL
002		015	055	000		DCC
002	300	302	355	002		JFZ ZEROFL

002 002 002 002 002	366 367 370	$\begin{array}{c} 257\\ 276 \end{array}$		002		LHLD 356 003 XRA CPM JFZ AHEAD3 LXH 000 005
002	376	021	360	003	AHEAD3,	LXD 360 003
003 003 003 003 003 003 003	002 003 006 007 010	176 247 312 022 023 043 303	014 001		BUFFMR,	LAM NDA JTZ NEXT STAD INXD INXH JMP BUFFMR
003 003 003 003 003 003 003	015 020 021 024 027	302	347 350	002	NEXT,	INXH SHLD 356 003 DCB JFZ NXWORD LXH 350 003 LMI 000 RET
	033 034 035 040				MSG,	LAM NDA RTZ CAL PRINT INXH JMP MSG

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